

Hardware Accelerated Next Generation Sequencing Analysis

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Workshop on Secure and Adaptive Computing from Genomics to Healthcare
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Sequencing data is increasing in volume and diversity



Illumina Genome
Analyzer, 2005
1 Gbases per day

Illumina NovaSeq X,
2025
4 Tbases per day



**4000x increase in
sequencing machine
throughput**

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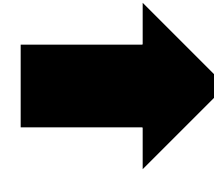
Illumina NovaSeq X,
2025
4 Tbases per day

4000x increase in
sequencing machine
throughput



illumina[®]

Read length: 100-350bp
Per base inaccuracy: 0.1%



**Oxford
NANOPORE
Technologies**

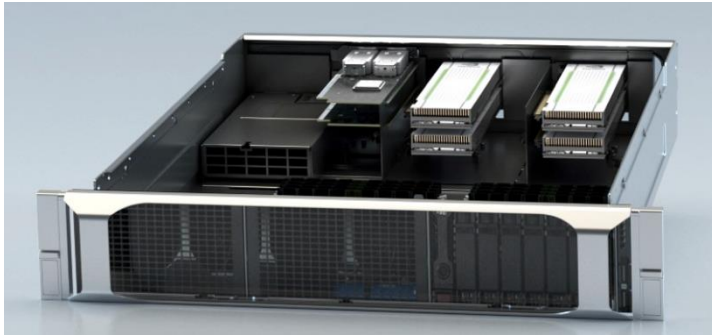
Read length: 1kb-1Mbps
Per base inaccuracy: 1-15%

1000x increase in
sequencing fragment
length



10 - 100x increase in
sequencing error rate

GPUs and custom hardware are being used for sequence analysis



GPU



Microsoft Genomics and
Stanford University



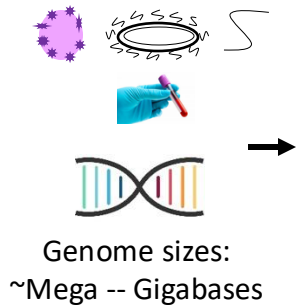
illumina® DRAGEN
Bio-IT Platform



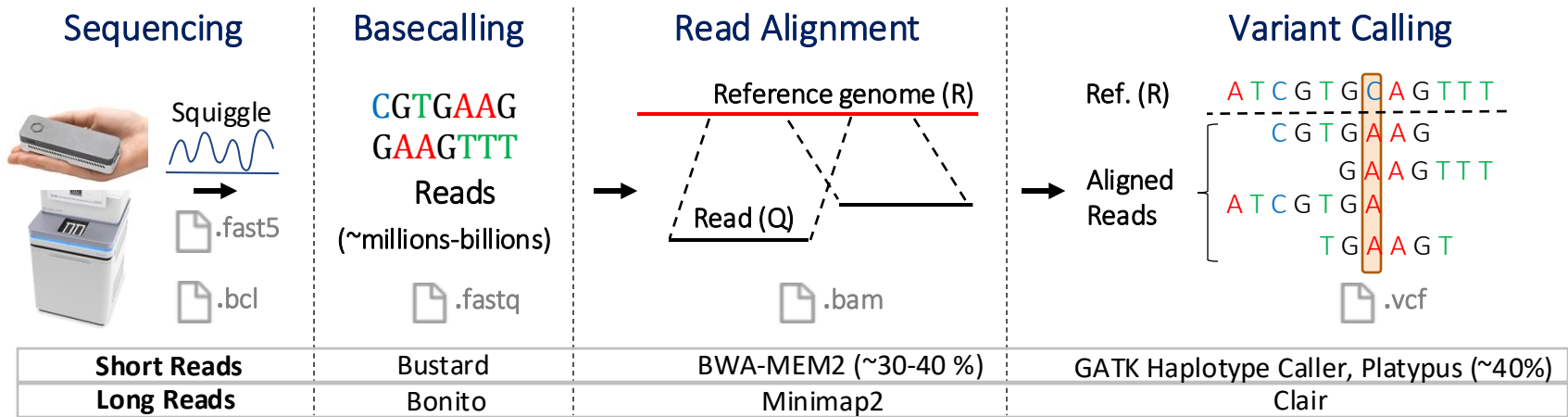
FPGA

Modern Sequence Analysis Pipelines

(a) Reference-Guided Assembly

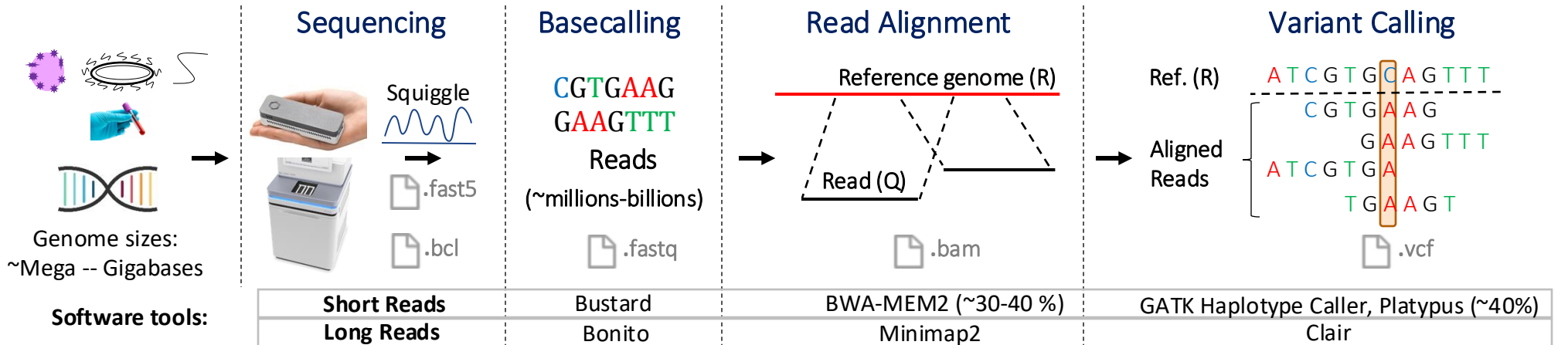


Software tools:



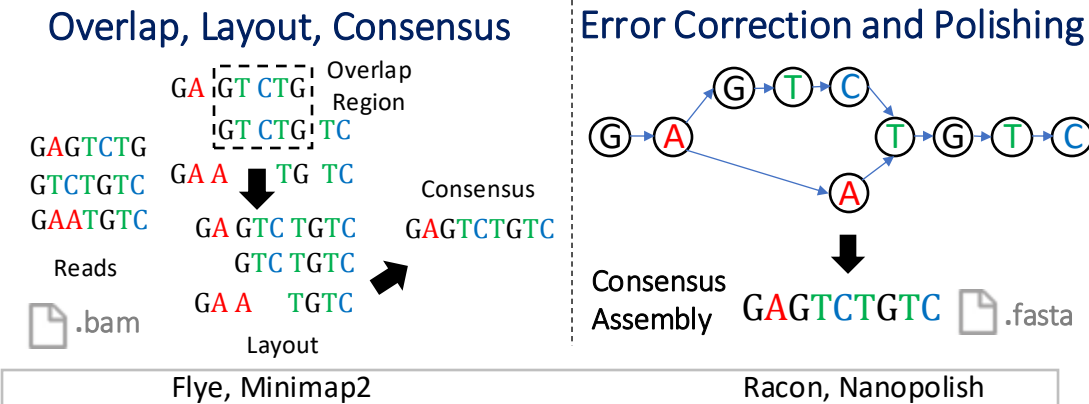
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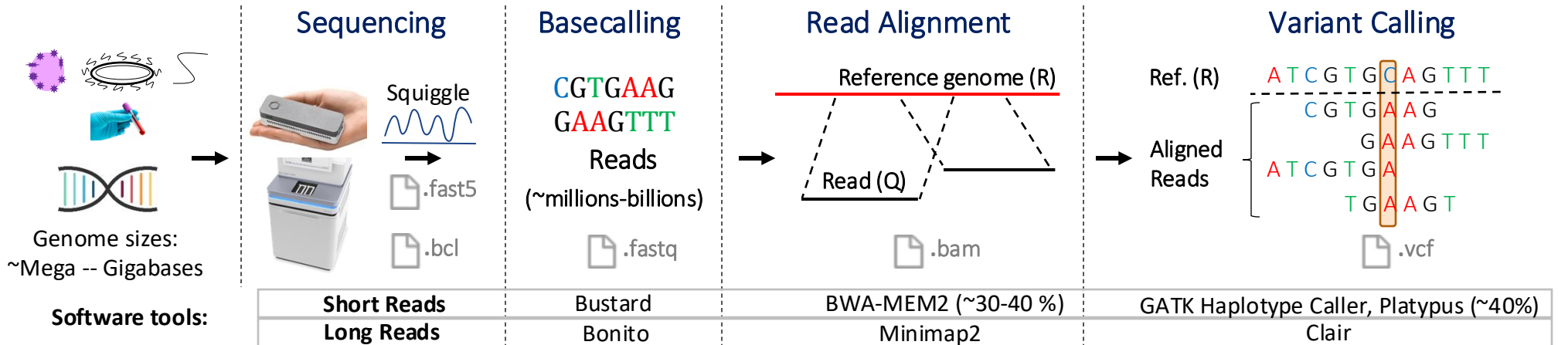
(b) De-Novo Assembly

Long Reads

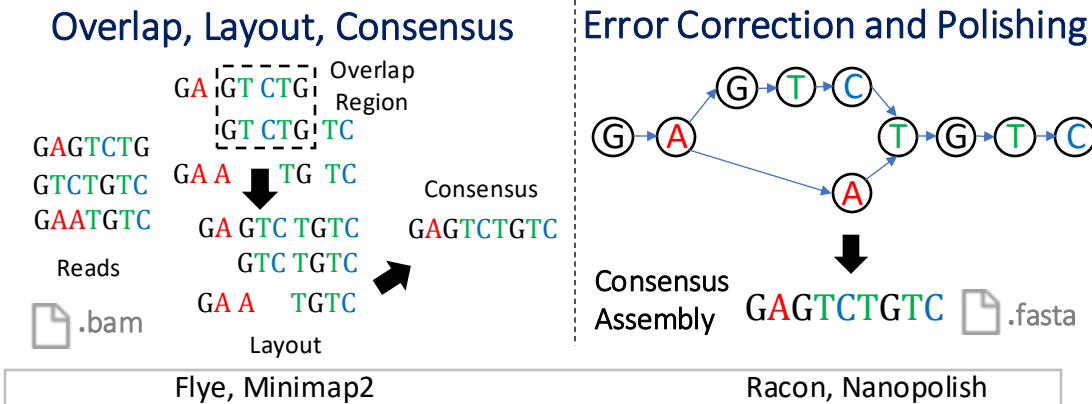


Modern Sequence Analysis Pipelines

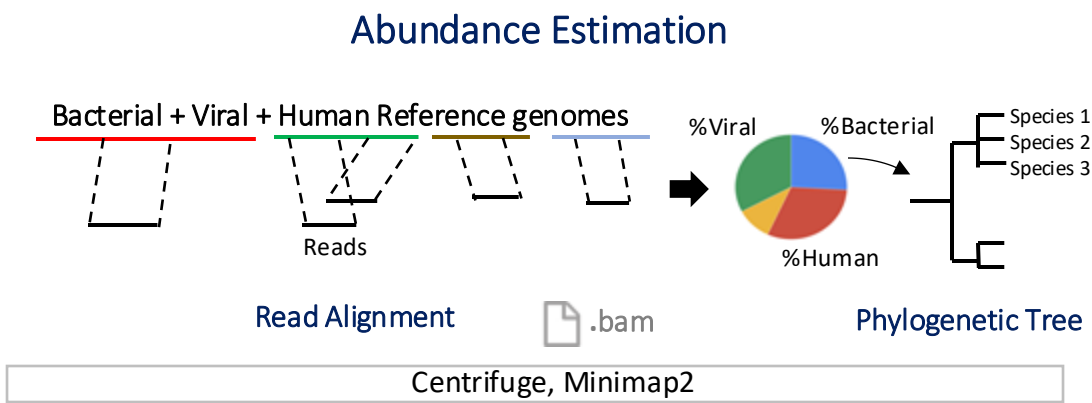
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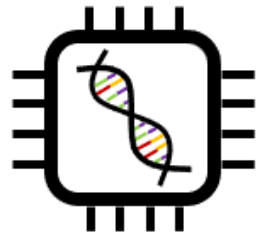


(b) De-Novo Assembly



(c) Metagenomics Classification





GenomicsBench

<https://genomicsbench.eecs.umich.edu>

- Computationally intensive kernels drawn from well maintained software tools
- Covers the major steps of modern sequence analysis pipelines
- Includes both short and long read analysis algorithms
- Small/large input datasets

Dynamic programming algorithms are common

Benchmark	Input Datatype	Applications	Chosen Tool	% Time Spent in Tool (single-thread)	Parallelism Motif
fmi	Short reads	<u>Read Alignment</u> Metagenomics Classification	BWA-MEM2	38%	Tree Traversal
bsw	Short reads	<u>Read Alignment</u> De-Novo Assembly	BWA-MEM2	31%	Dynamic Programming
dbg	Short reads	<u>Variant Calling</u> De-Novo Assembly	Platypus	65%	Graph Construction Hash Table
phmm	Short reads	<u>Variant Calling</u> Error Correction	GATK Haplotype Caller	70%	Dynamic Programming
chain	Long reads	<u>De-Novo Assembly</u> <u>Read Alignment</u>	Minimap2	47.4 %	Dynamic Programming (1D)
spoa	Long reads	Error Correction	Racon	75 %	Dynamic Programming Graph Construction
abea	Long reads	<u>Basecalling</u> <u>Variant Calling</u>	Nanopolish	71.4%	Dynamic Programming
grm	NA	Population Genomics	PLINK2	92.8 %	Dense Matrix Multiplication
nn-base	Long reads	Basecalling	Bonito	95 %	FP Matrix Multiplication
nn-variant	Long reads	Variant Calling	Clair	57.2 %	FP Matrix Multiplication

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Dynamic programming algorithms are common

- 5 dynamic programming kernels in GenomicsBench
- Contribute 30-75% execution time in their respective software tools

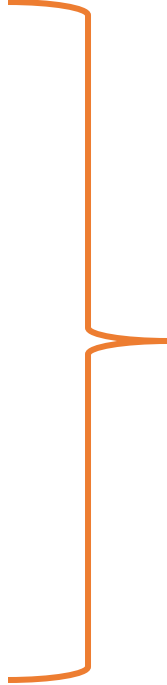
Dynamic programming algorithms are common

- 5 dynamic programming kernels in GenomicsBench
- Contribute 30-75% execution time in their respective software tools
 - Chaining
 - Banded Smith-Waterman
 - Pairwise Hidden Markov Model
 - Partial Order Alignment
 - Adaptive Banded Event Alignment

Dynamic programming algorithms are common

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Diverse compute requirements

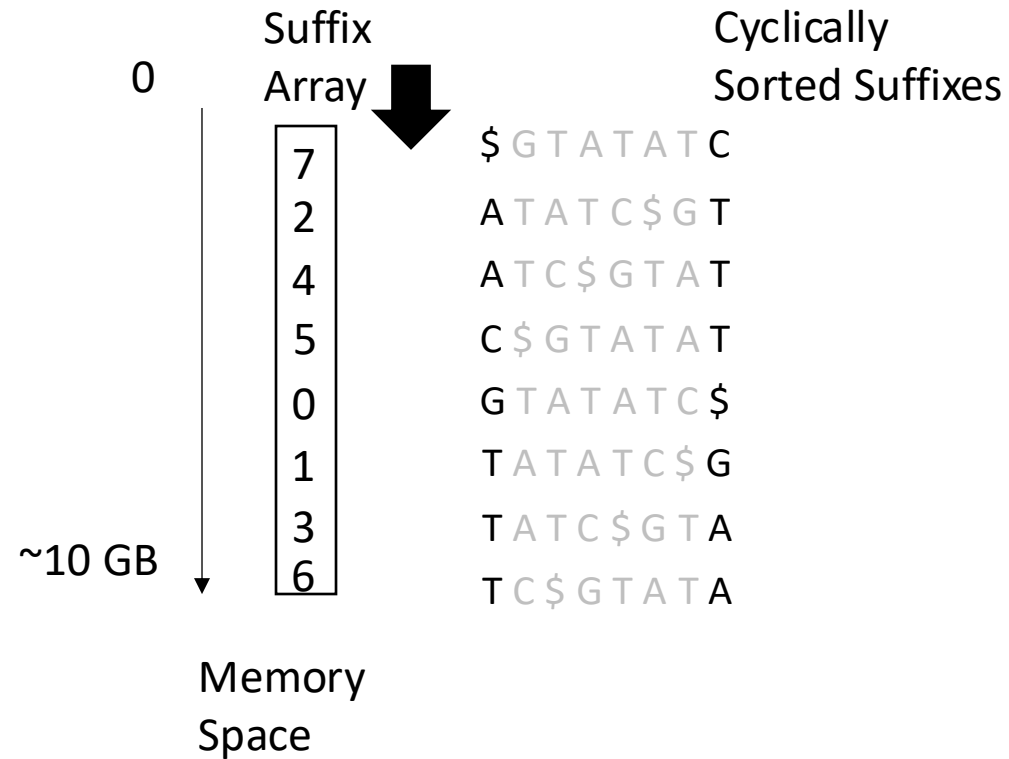
Data dependencies: 1-D, 2-D

Inputs: Sequence, Graph

Computation type: Integer, Floating point

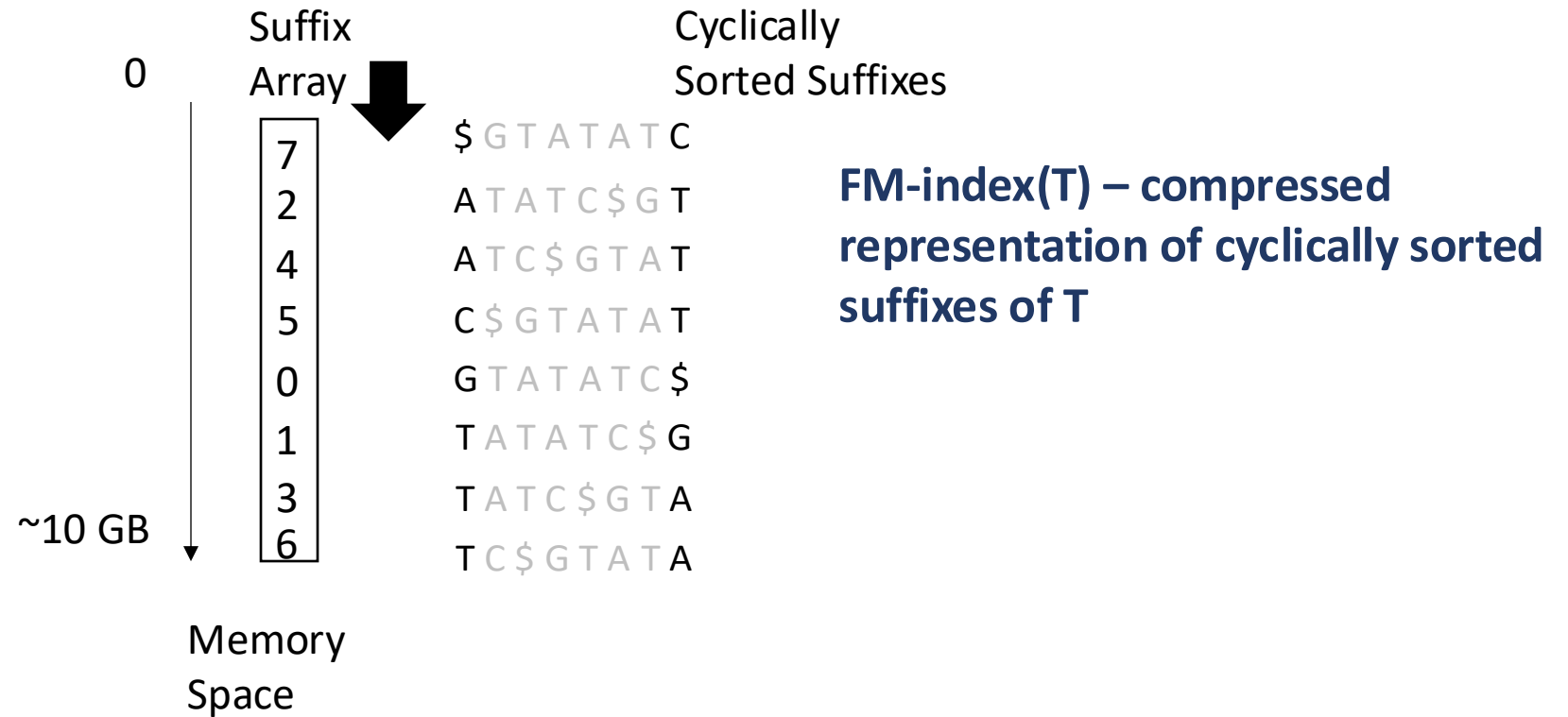
Memory-Intensive Index Lookup Benchmarks: e.g., FM-index search

Reference : **G T A T A T C \$**



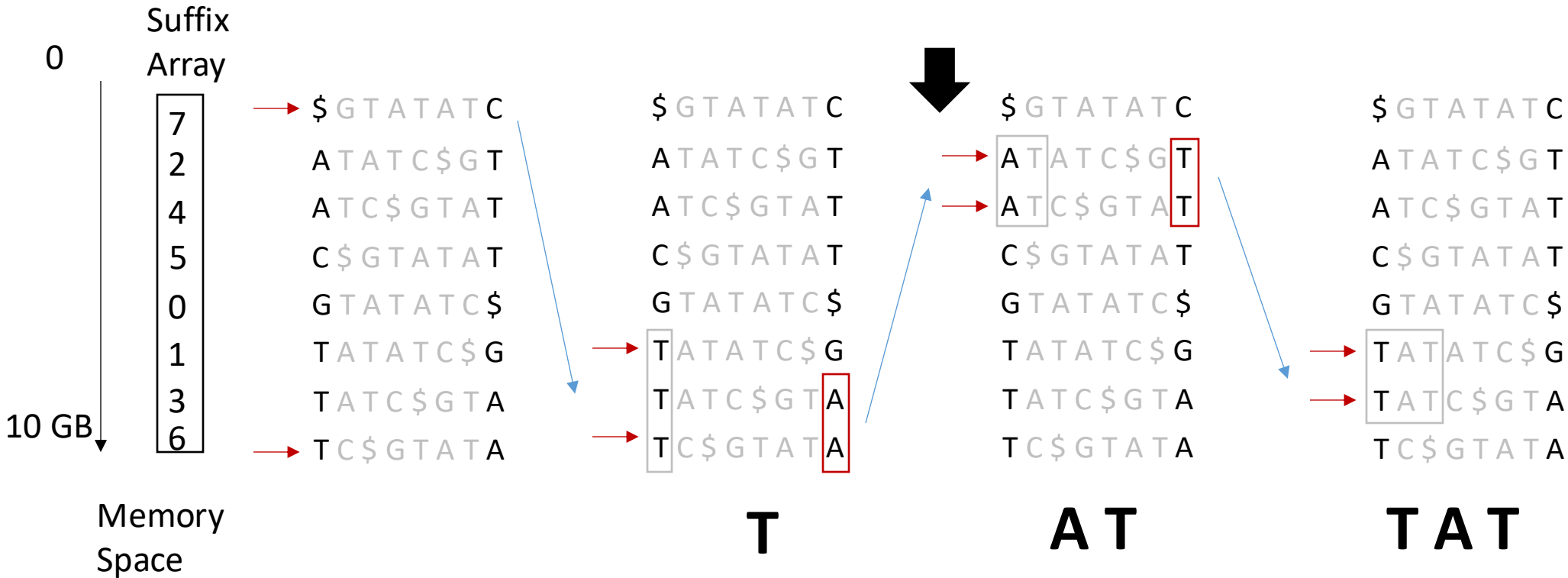
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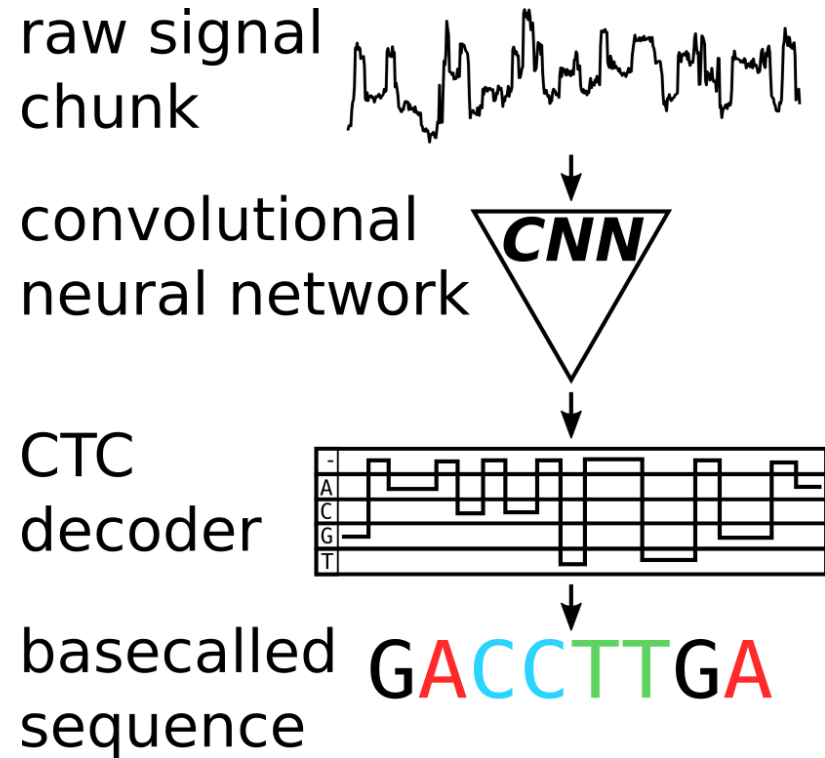
Irregular memory accesses in FM-index search

Reference : G T A T A T C \$



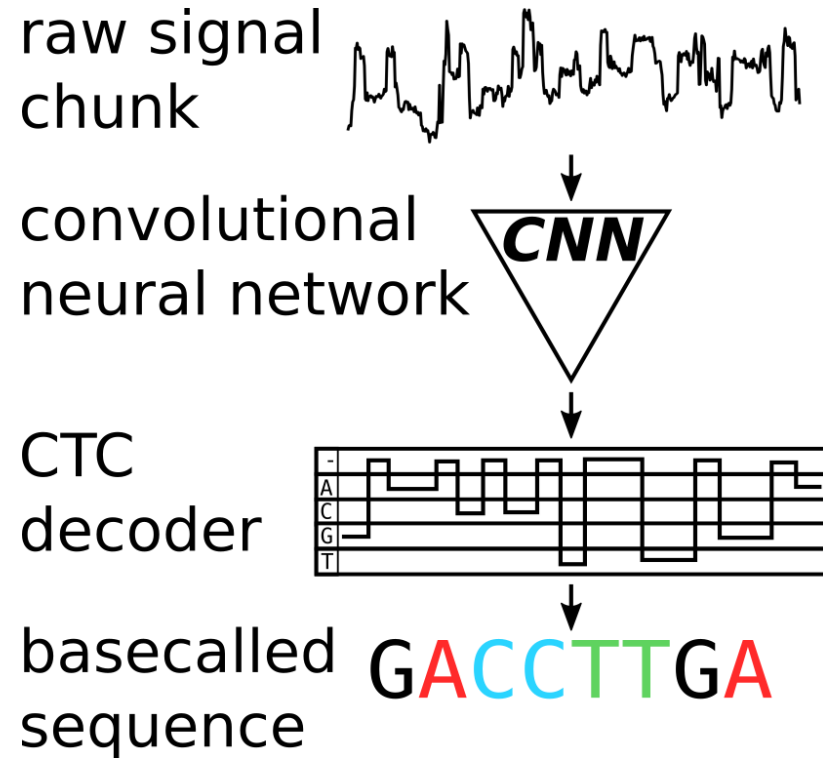
Read : T A T

Machine learning algorithms are also gaining traction

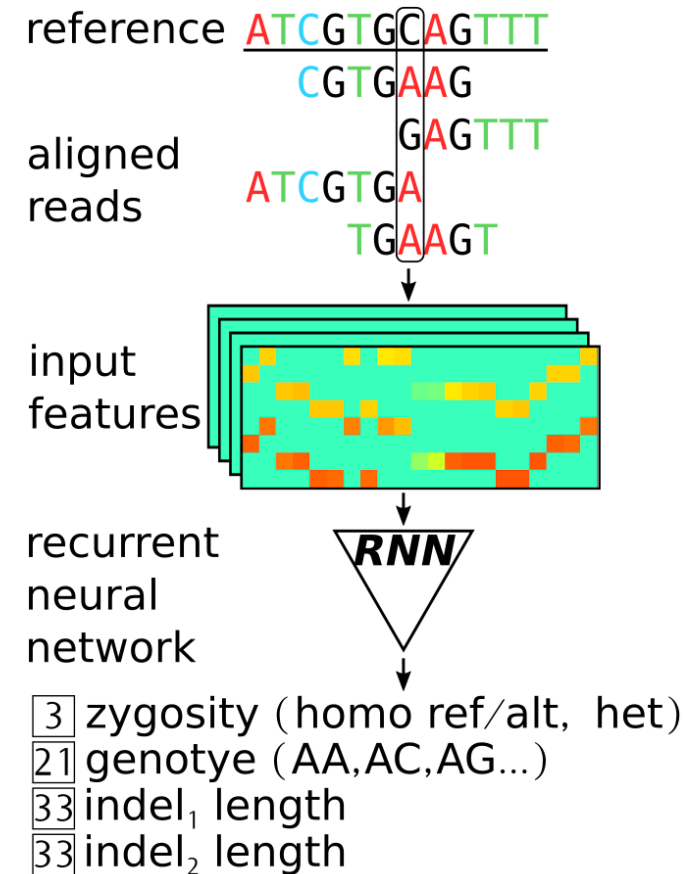


nn-base (Basecalling)

Machine learning algorithms are also gaining traction



nn-base (Basecalling)



nn-variant (Variant calling)

Abundant data parallelism in sequence analysis applications

Kernel	Parallelism granularity	Source of parallelism
FM-index search (fmi)	Read	# OCC Table Lookups
Banded Smith-Waterman (bsw)	Seed	# Cell Updates
Pairwise Hidden Markov Model (phmm)	Genome Region	# Cell Updates
Chaining (chain)	Read	# Input Anchors
Partial Order Alignment (poa)	Read Chunk	# Cell Updates

Abundant data parallelism available

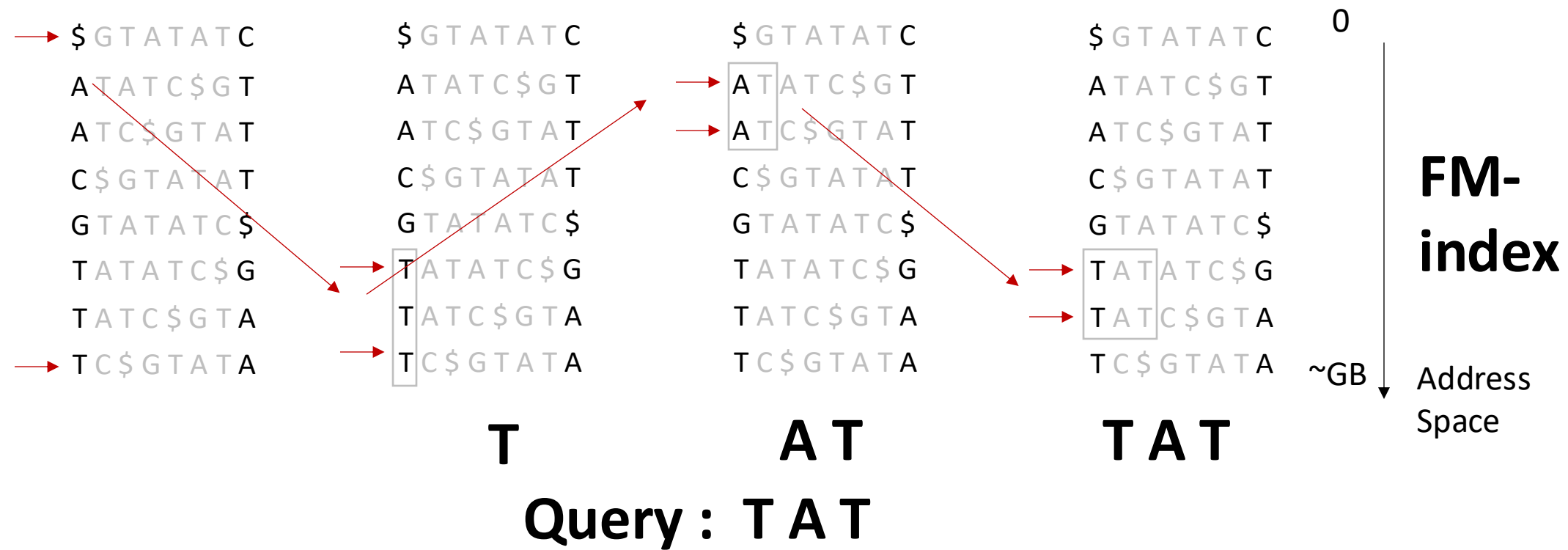
~billions of *independent* reads

~millions of *independent* genome regions

Case Study – Accelerating FM-index Search

FMD-index seeding trades off memory bandwidth for space

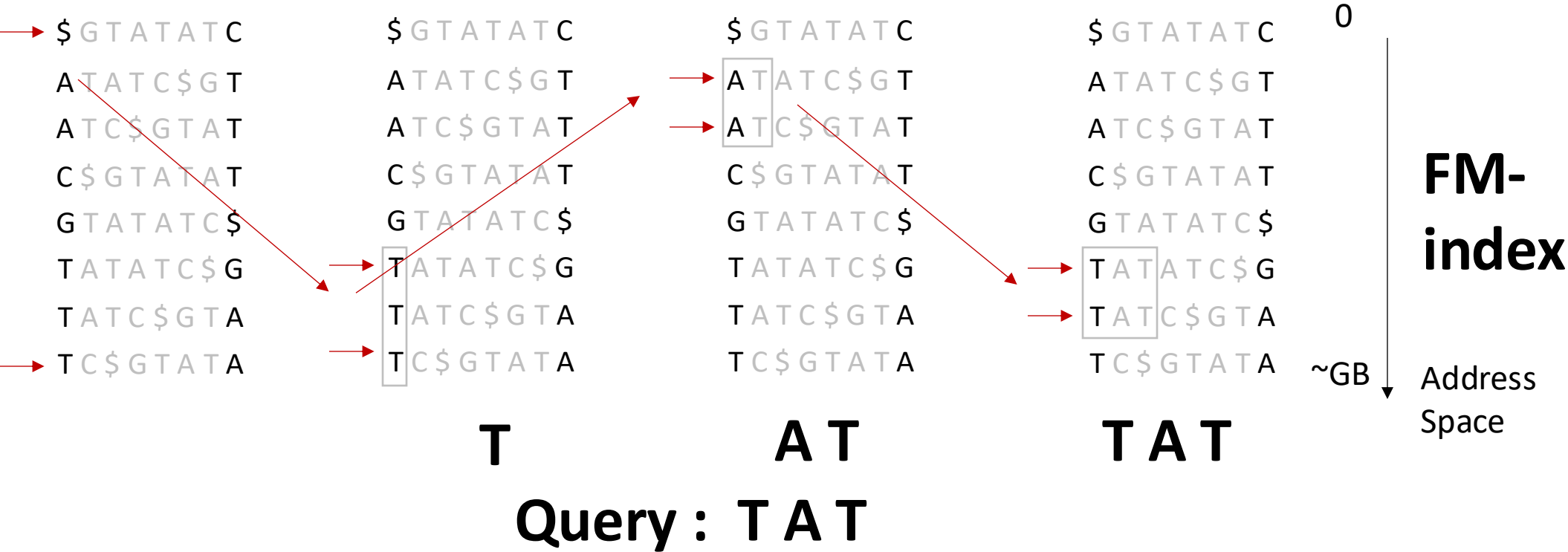
! Single Character Lookup



FMD-index seeding trades off memory bandwidth for space

! Single Character Lookup

Irregular
memory
accesses
!



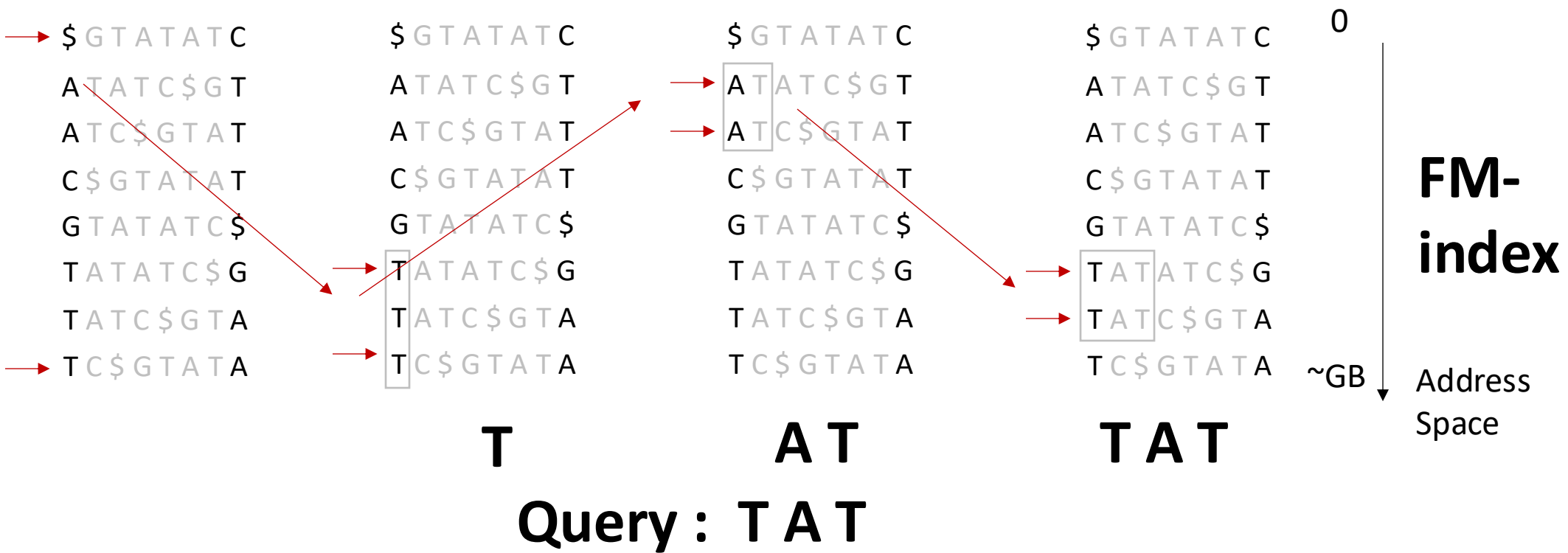
FMD-index seeding trades off memory bandwidth for space

BWA-MEM : 102 kB / read

BWA-MEM2 : 61 kB / read

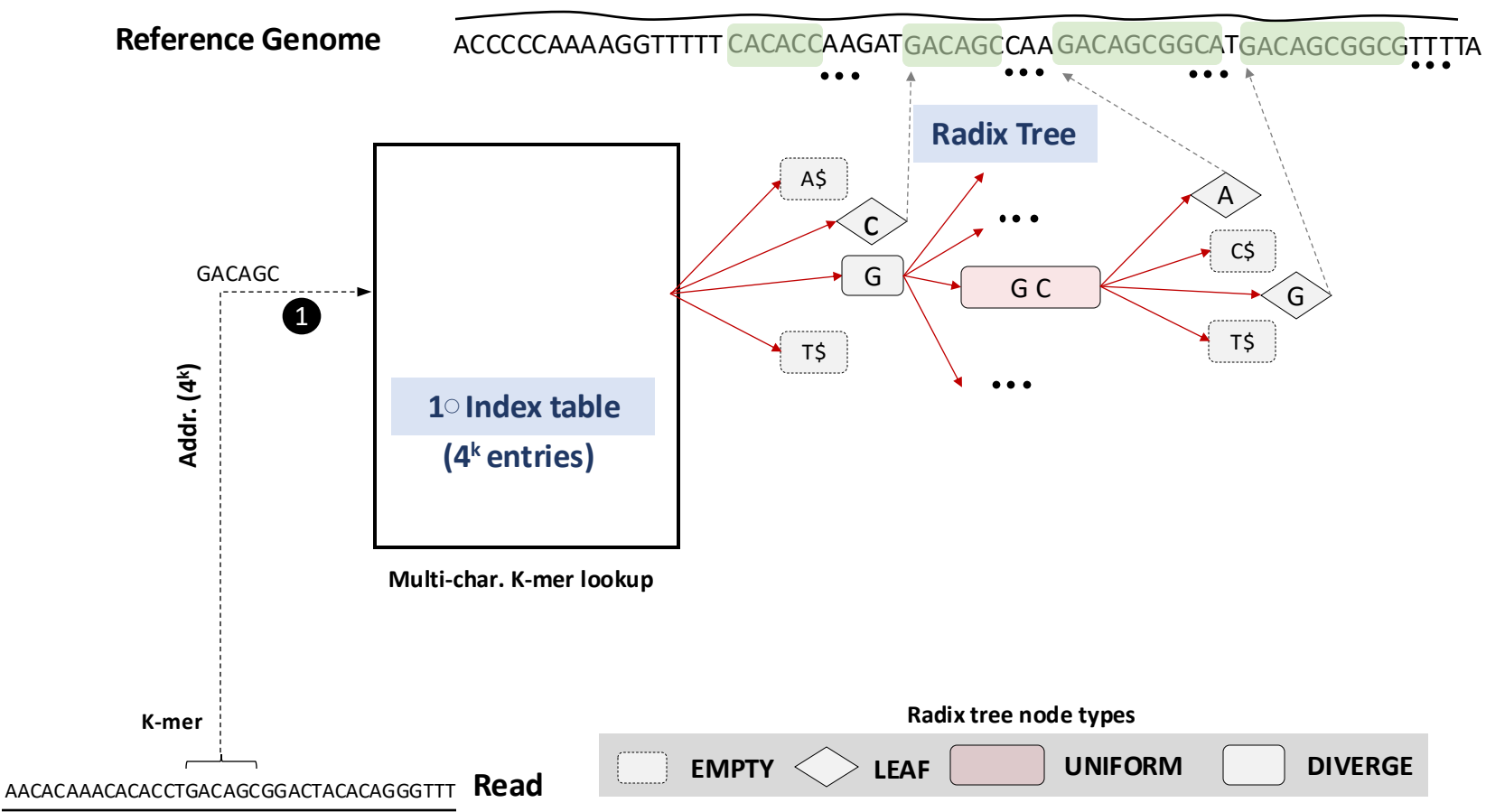
! Single Character Lookup

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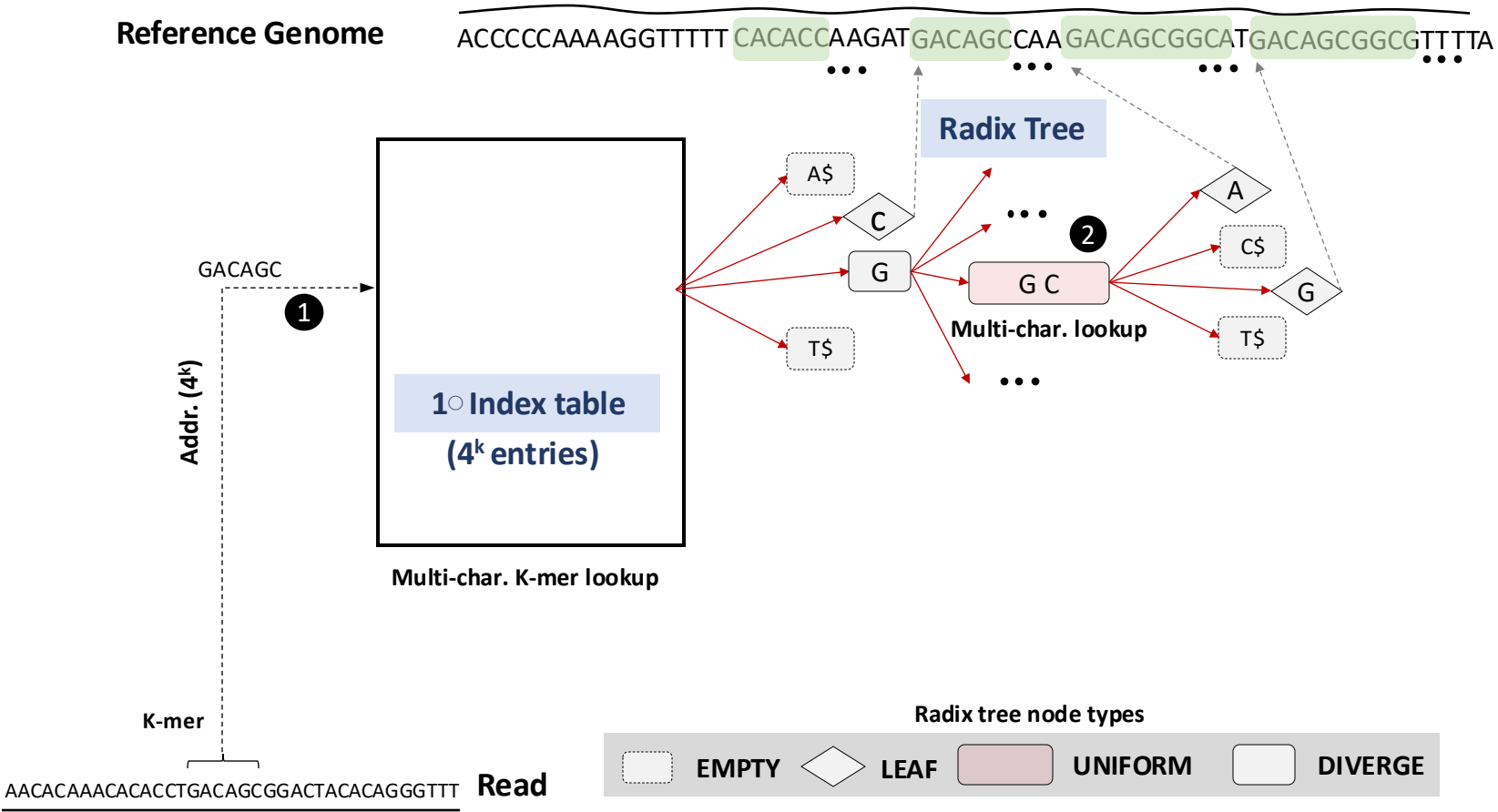
Multi-character lookup trading space for reducing memory bandwidth

ERT: Multi-level index table with a space-optimized radix tree



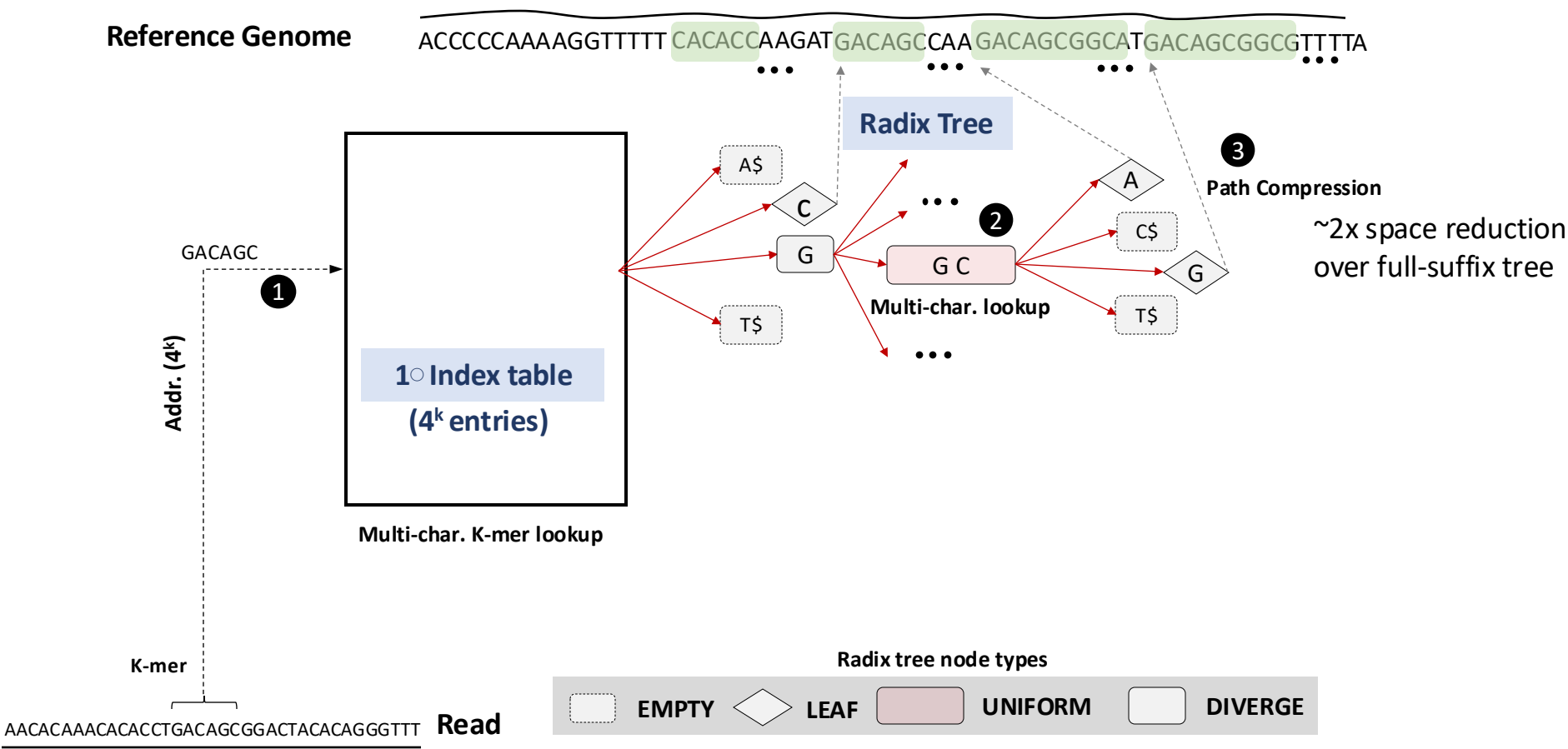
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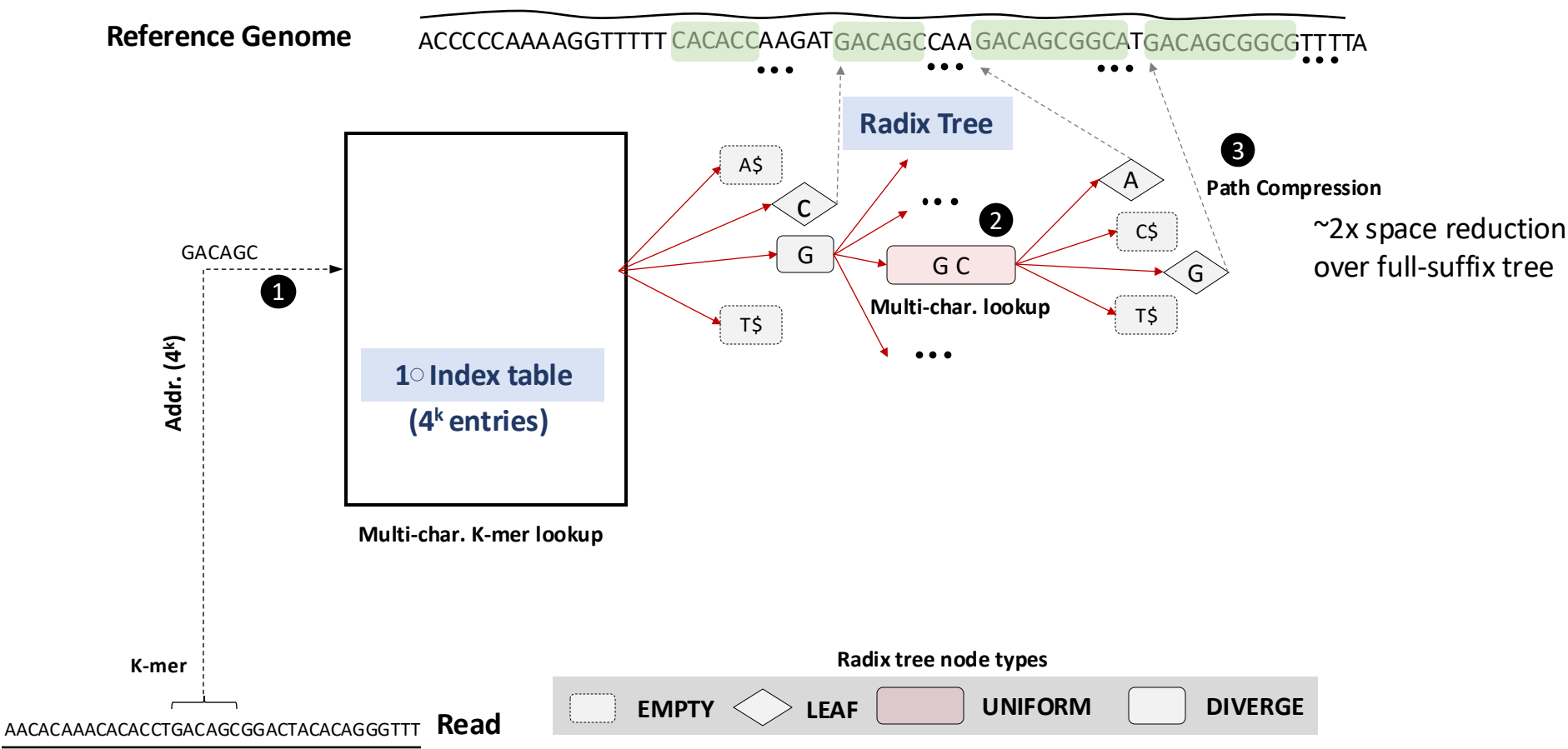
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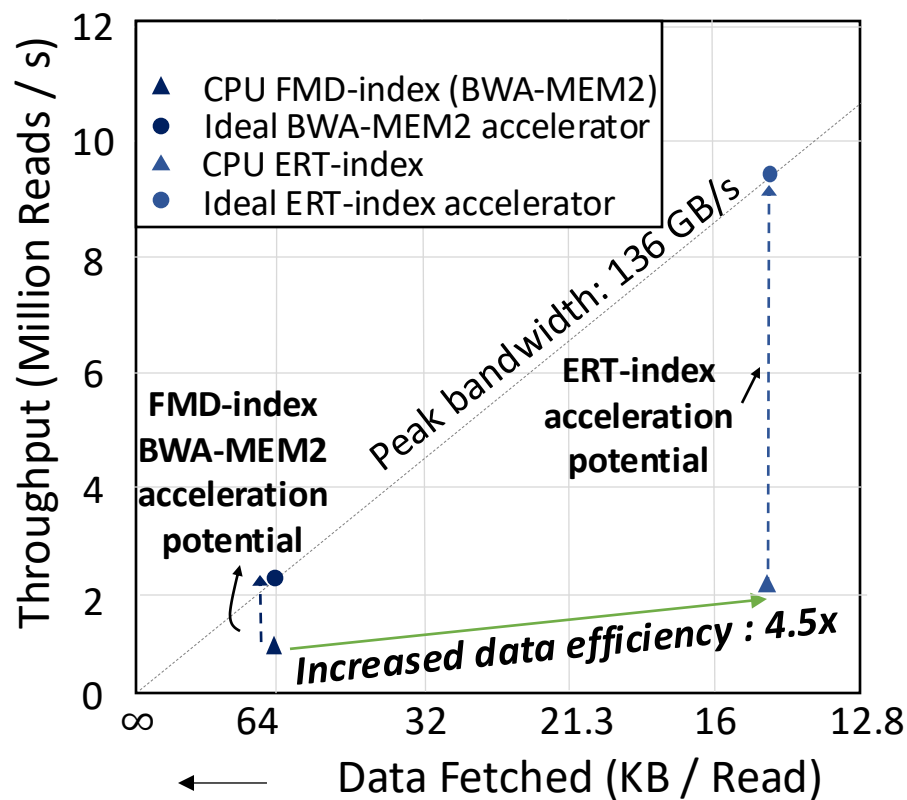
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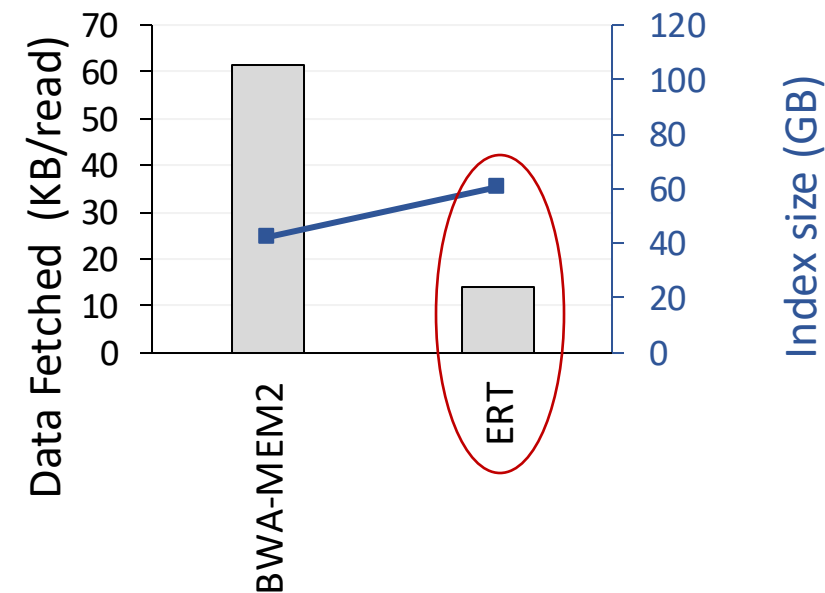


Also, present is a 2nd level index table for 15-mers having ≥ 256 occurrences

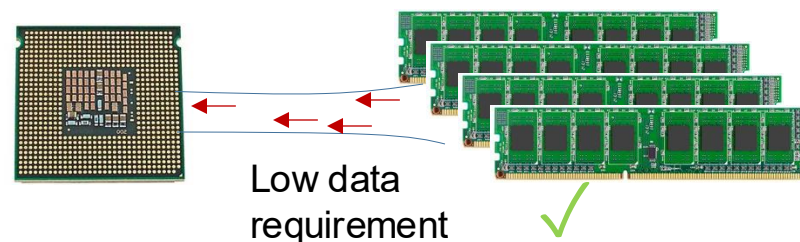
Instead trade-off memory capacity for reducing memory bandwidth



■ Data fetched (KB/Read) —■ Index Size (GB)



Enumerated
Radix Tree
(ERT)



Low data
requirement

ERT
~60 GB
human

Dynamic programming for sequence similarity estimation

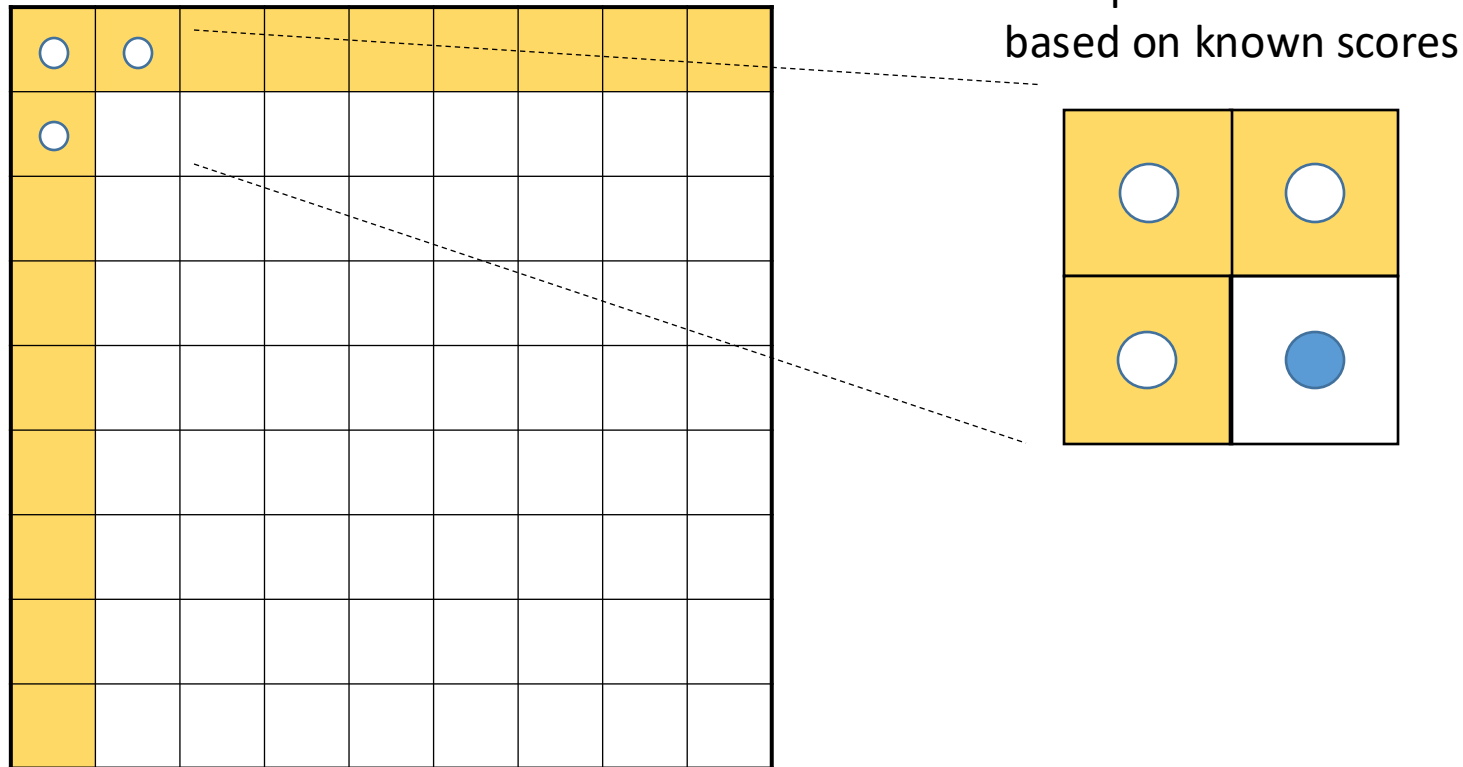
- Task: calculate the scores in the entire table.

Dynamic programming for sequence similarity estimation

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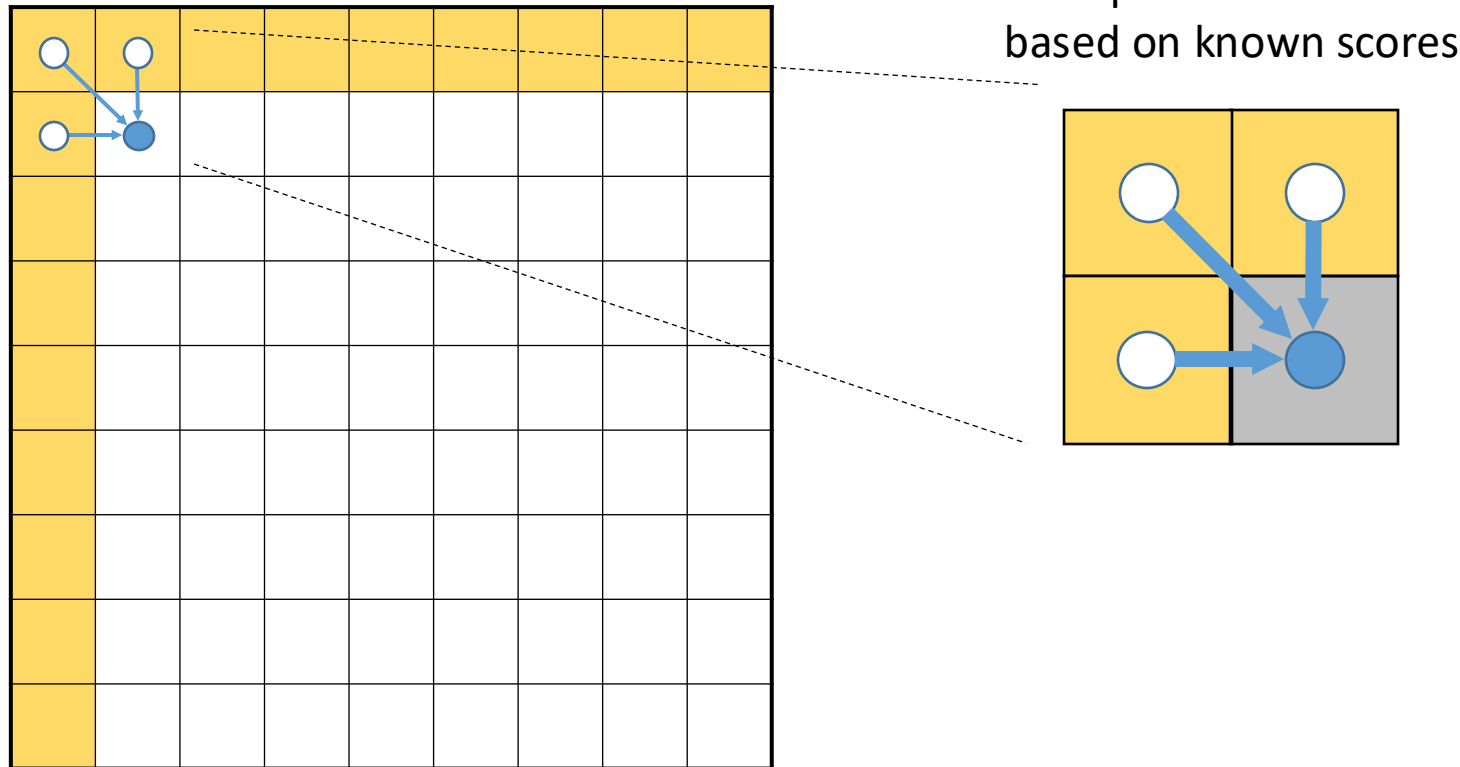
Dynamic programming for sequence similarity estimation

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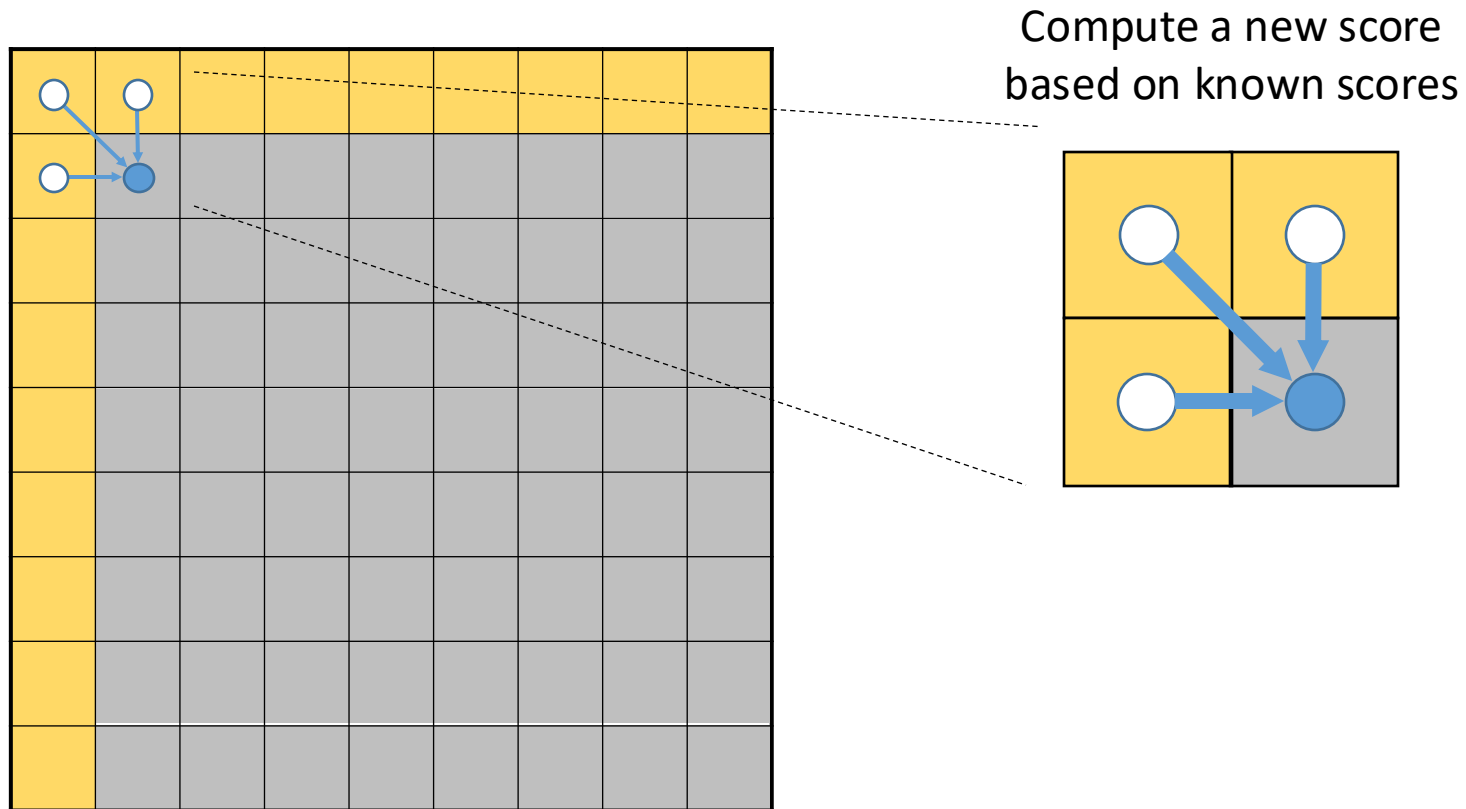
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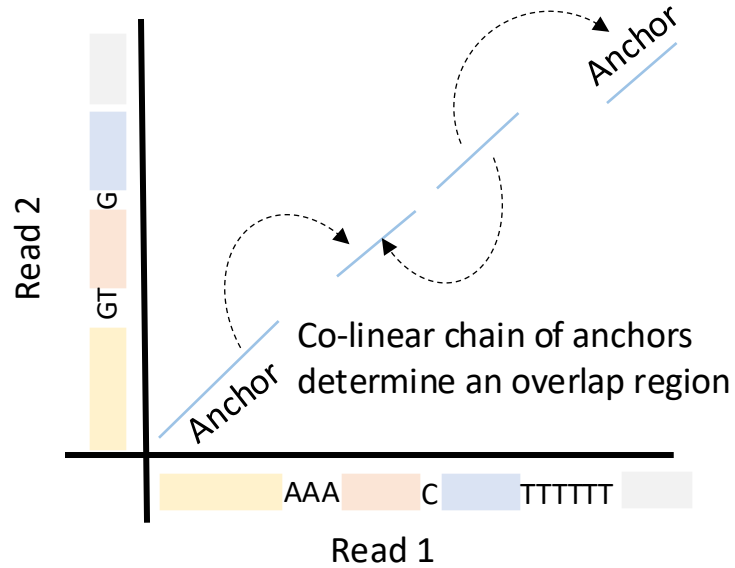


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1D dynamic programming



$$\boxed{score(i)} = \max \left\{ \max_{i > j \geq 1} \left\{ \boxed{score(j)} + \boxed{\alpha(j, i)} - \boxed{\beta(j, i)} \right\}, \boxed{w_i} \right\}$$

Score for anchor i

Score for anchor j

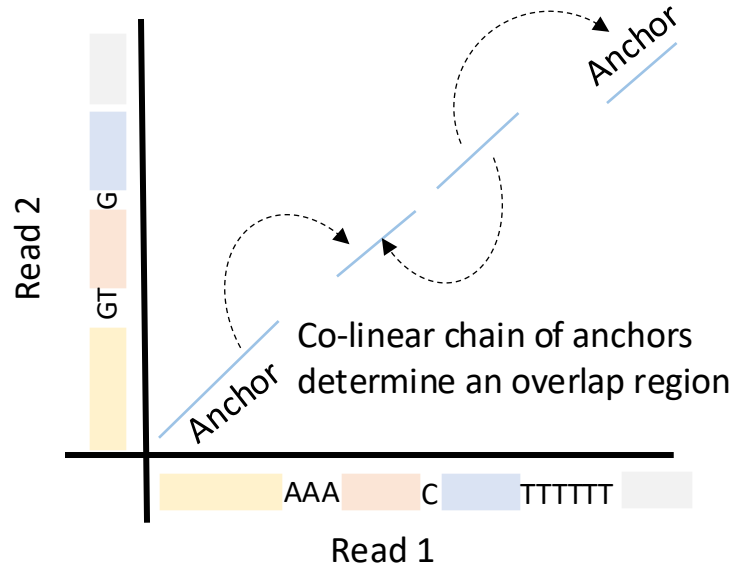
matching bases

gap length

anchor length

Chaining

1D dynamic programming



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Score for anchor i

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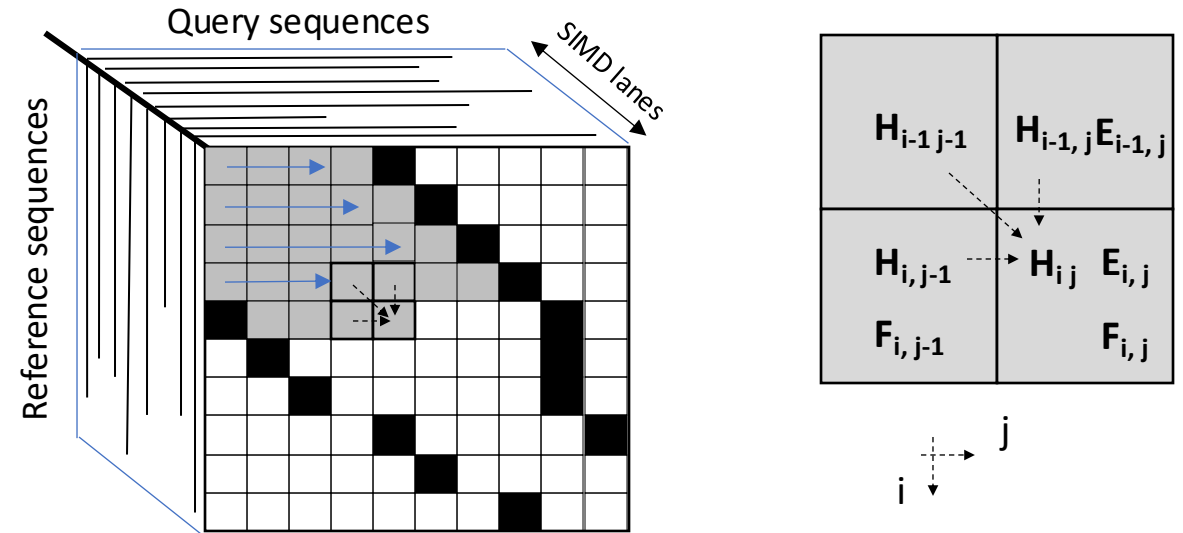
matching bases

gap length

anchor length

Chaining

2D dynamic programming



$$H_{ij} = \max \{ H_{i-1,j-1} + s(i, j), E_{ij}, F_{ij} \}$$

$$E_{i+1,j} = \max \{ H_{ij} - q, E_{ij} \} - e$$

$$F_{i,j+1} = \max \{ H_{ij} - q, F_{ij} \} - e$$

Affine gap scoring

Banded Smith-Waterman

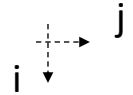
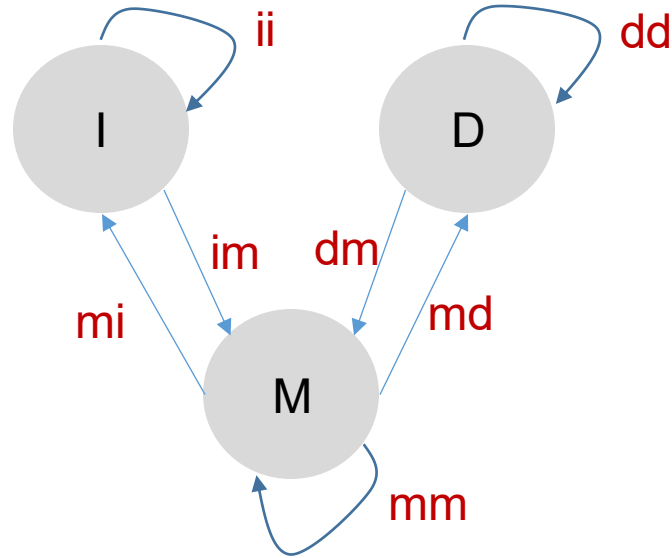
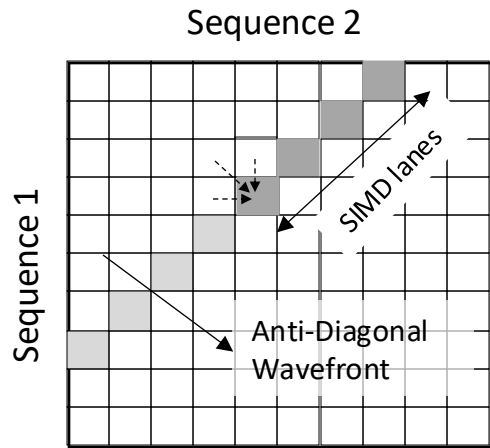
Dynamic programming benchmarks: Floating point 2D

Pairwise Hidden Markov Model



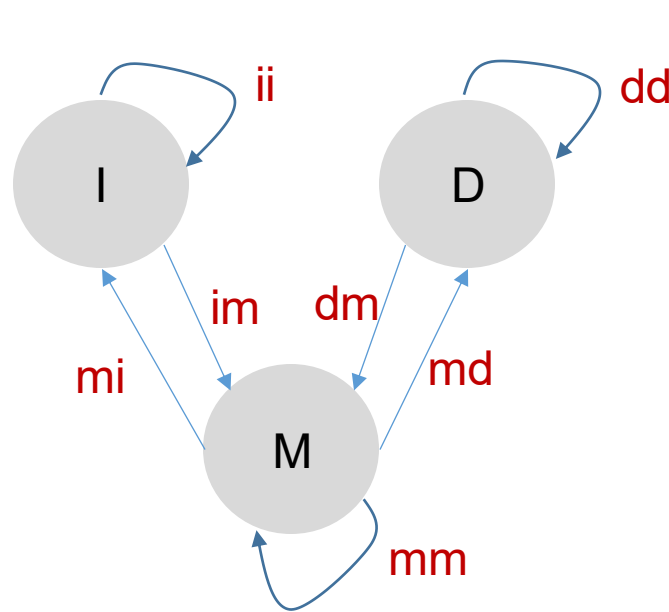
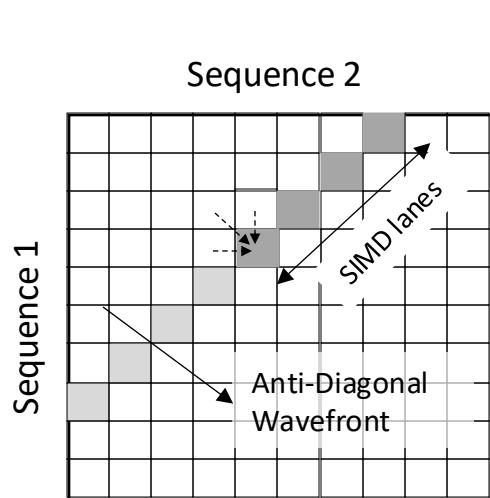
Dynamic programming benchmarks: Floating point 2D

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Dynamic programming benchmarks: Floating point 2D

Pairwise Hidden Markov Model



$$M(i,j) = (\begin{array}{l} mm \cdot M(i-1,j-1) \\ + im \cdot I(i-1,j-1) \\ + dm \cdot D(i-1,j-1) \end{array})$$

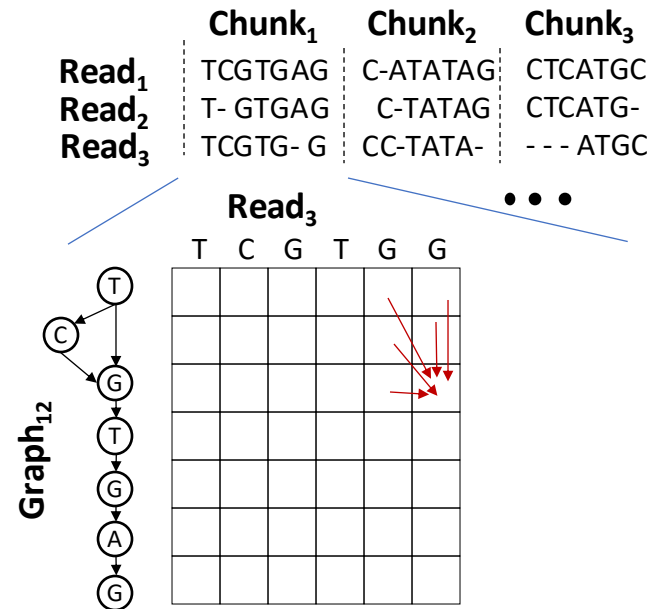
$$I(i,j) = mi \cdot (i-1,j) + ii \cdot I(i-1,j)$$

$$D(i,j) = md \cdot (i,j-1) + dd \cdot D(i,j-1)$$

Floating Point

Dynamic programming benchmarks: Sequence to Graph

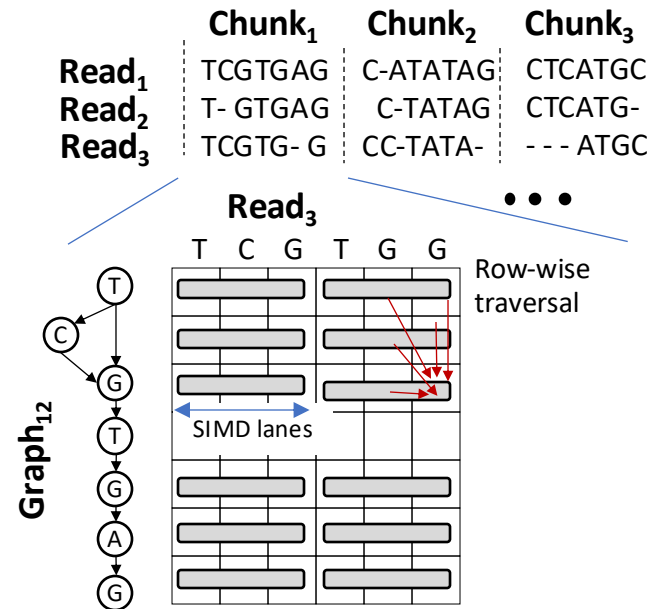
Partial Order Alignment



Sequence to Graph Alignment

Dynamic programming benchmarks: Sequence to Graph

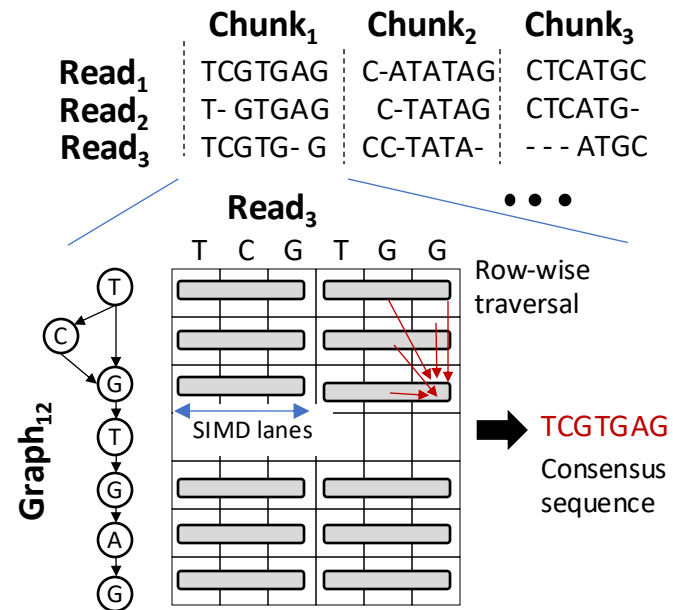
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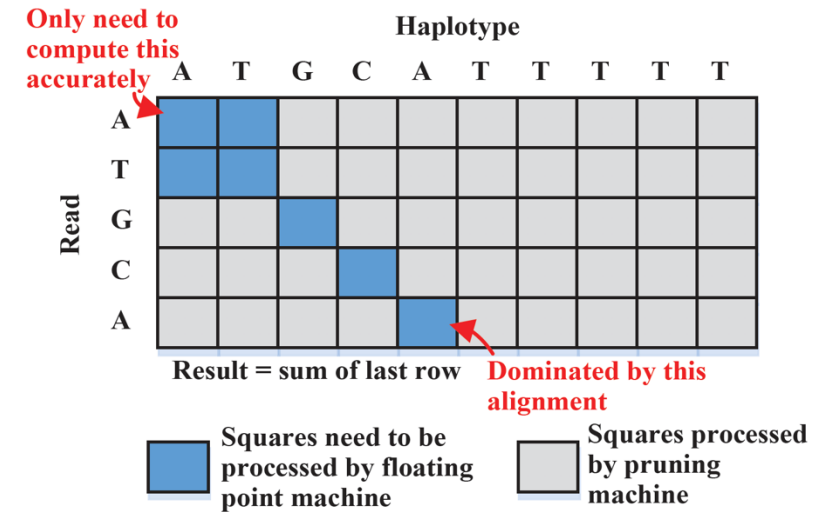


Sequence to Graph Alignment

Case Study – Accelerating Pairwise Hidden Markov Model (pairHMM)

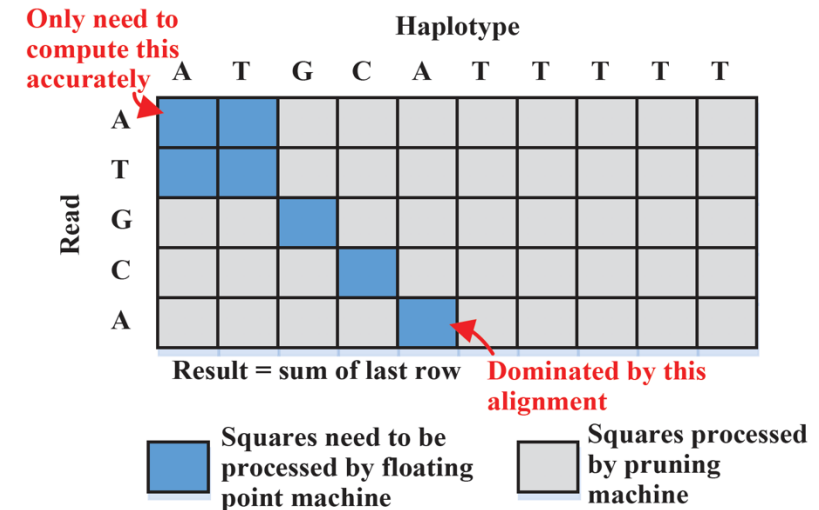
Pruning algorithm for pairHMM with same output guarantees

- Compute upper bound probability by replacing floating-point multiplication with log-domain addition and rounding up (high_{res})
- Prune matrix and identify promising paths
- Compute lower bound probability with single-precision floating-point only for small unpruned area (low_{res})
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Original algorithm

Downstream: Call the most likely mutation

	Genotype	Probability
G1	GG	0.1
G2	AA	0.2
G3	GA	0.1

Called

Pruned version

Downstream: Bound checks & call mutation

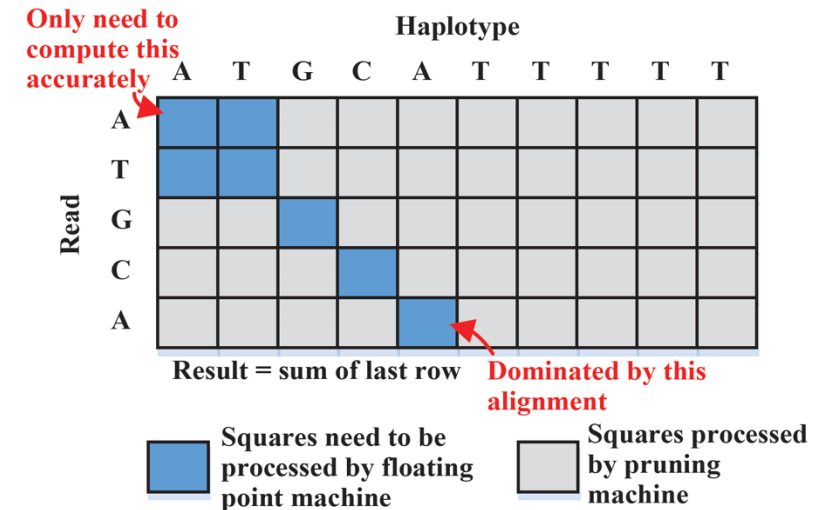
	Genotype	Probability
G1	GG	[0.05, 0.15]
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Called

Nonoverlapping bounds:
G2 is the correct call

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Called

Nonoverlapping bounds:
G2 is the correct call

Results:

43x fewer cells computed in precise floating point

8.3x higher throughput (GCUPS) than floating-point ASIC of the same area

Case Study – Generalized Dynamic Programming Acceleration

DPX Instruction Extension in NVIDIA Hopper GPU



DPX Instruction Extension in NVIDIA Hopper GPU

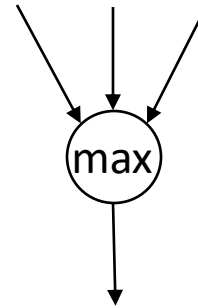
ISA Extension



DPX Instruction Extension in NVIDIA Hopper GPU



ISA Extension

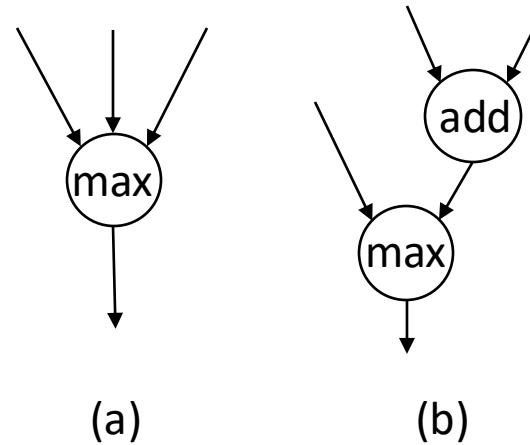


(a)

DPX Instruction Extension in NVIDIA Hopper GPU



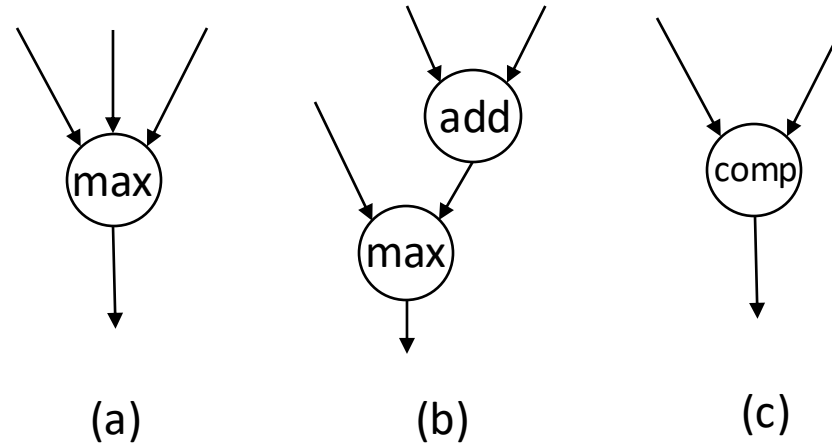
ISA Extension



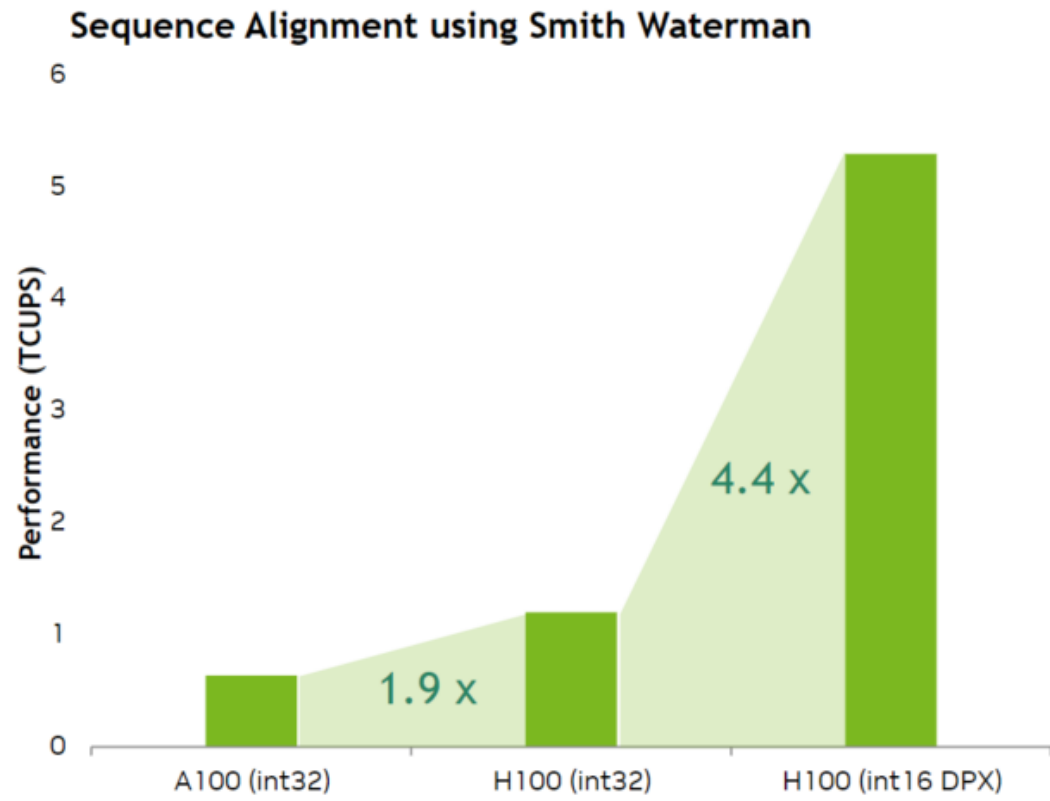
DPX Instruction Extension in NVIDIA Hopper GPU



ISA Extension

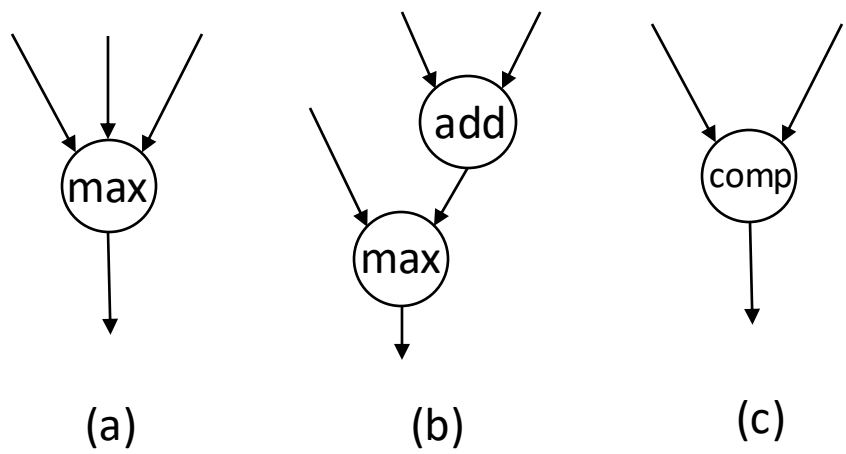


DPX Instruction Extension in NVIDIA Hopper GPU



Affine gap alignment (score calculation)
Weights used from BWA.
Data: [HG002 \(NA24385\)](#) paired-end protocol using Illumina Sequencers.

ISA Extension

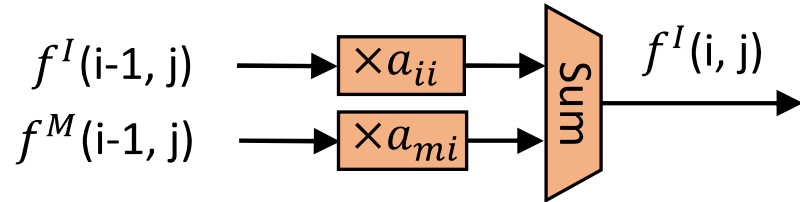


Diversity of dynamic programming in genomic kernels

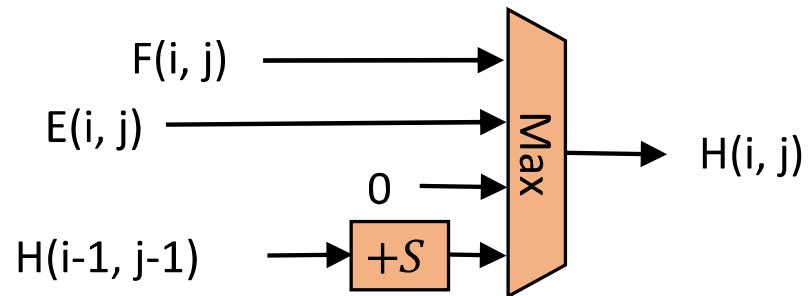
Kernel	Application	Dimension and Size	Dependency	Data Type
bsw	Read Alignment	2D $\sim 120 \times 60$	Last 2 Wave-fronts	Int 8/16
pairHMM	Variant Calling	2D $\sim 100 \times 60$	Last 2 Wave-fronts	Floating Point
poa	Error Correction	2D $\sim 1000 \times 500$	Graph Structure	Int 32
chain	Read Alignment	1D ~ 20000	Last N (~ 25) Anchors	Int 32

Different objective functions computed in each cell

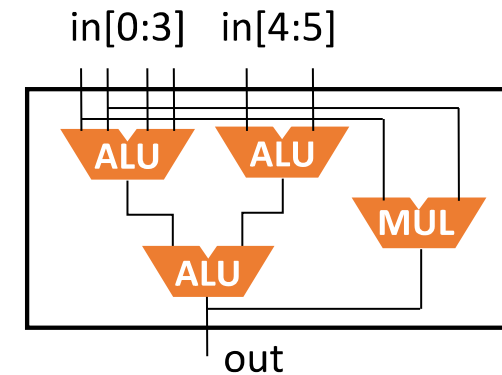
GenDP: Support for intra-cell objective functions



(a) Reduction Data-Flow in PairHMM

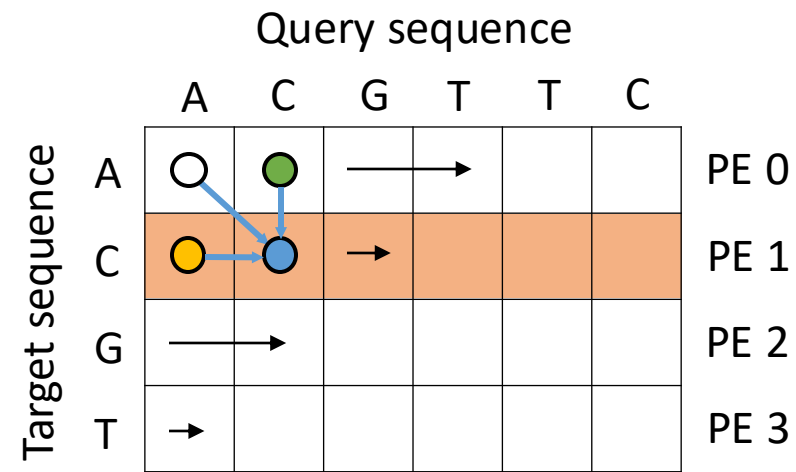


(b) Reduction Data-Flow in BSW

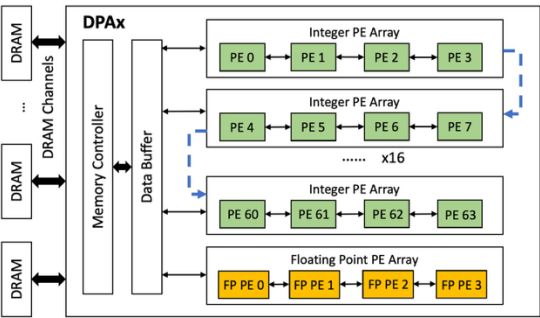


Compute unit – 2-level reduction tree

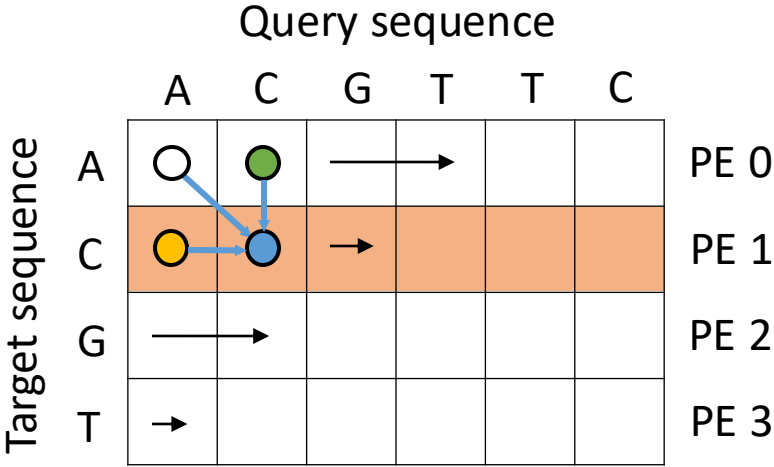
GenDP: Inter-cell dependencies and multi-precision support



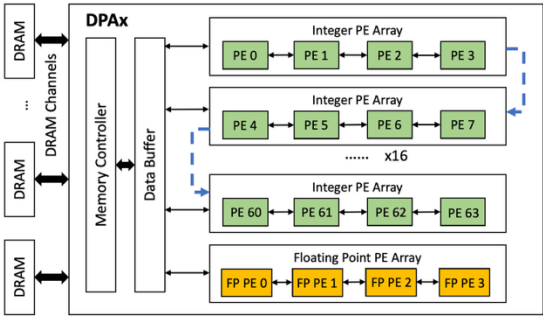
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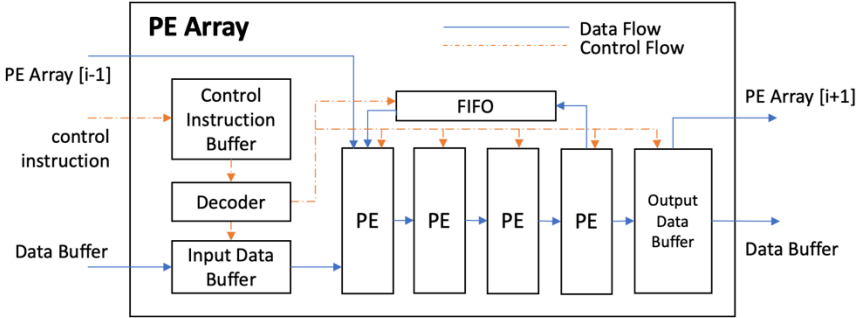
Integer and floating point PE array



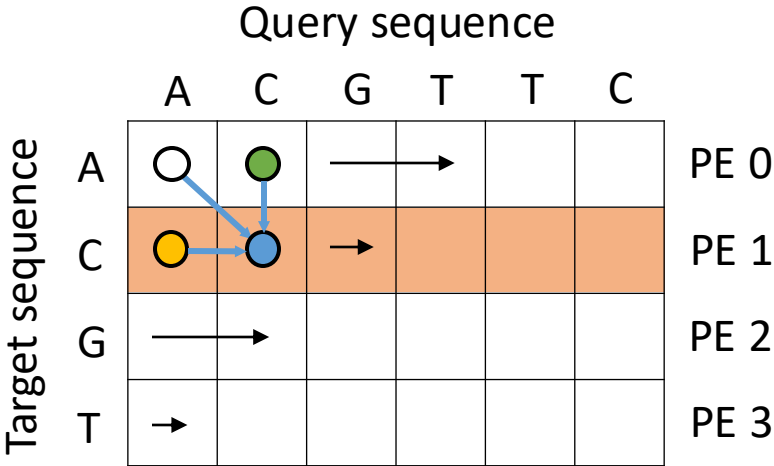
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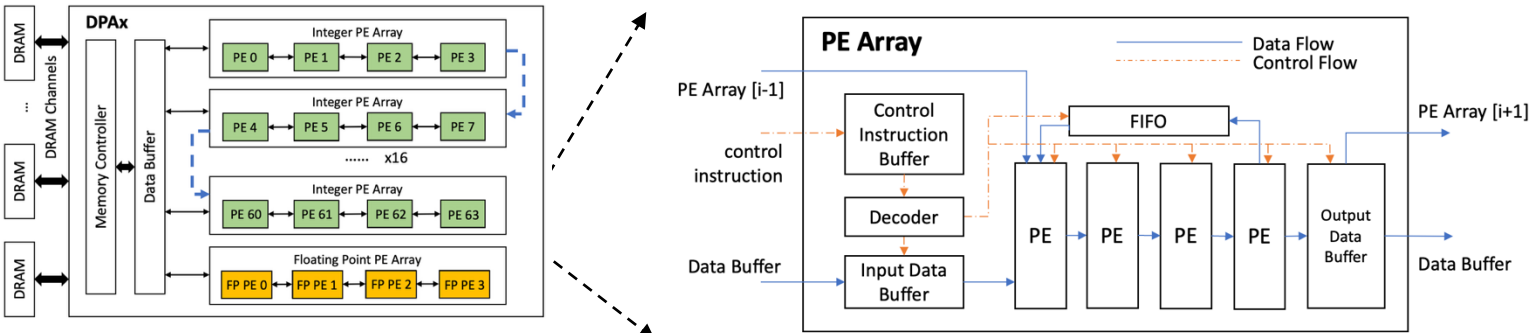
Integer and floating point PE array



Local dependency - 1-Dimension systolic PE array with FIFO

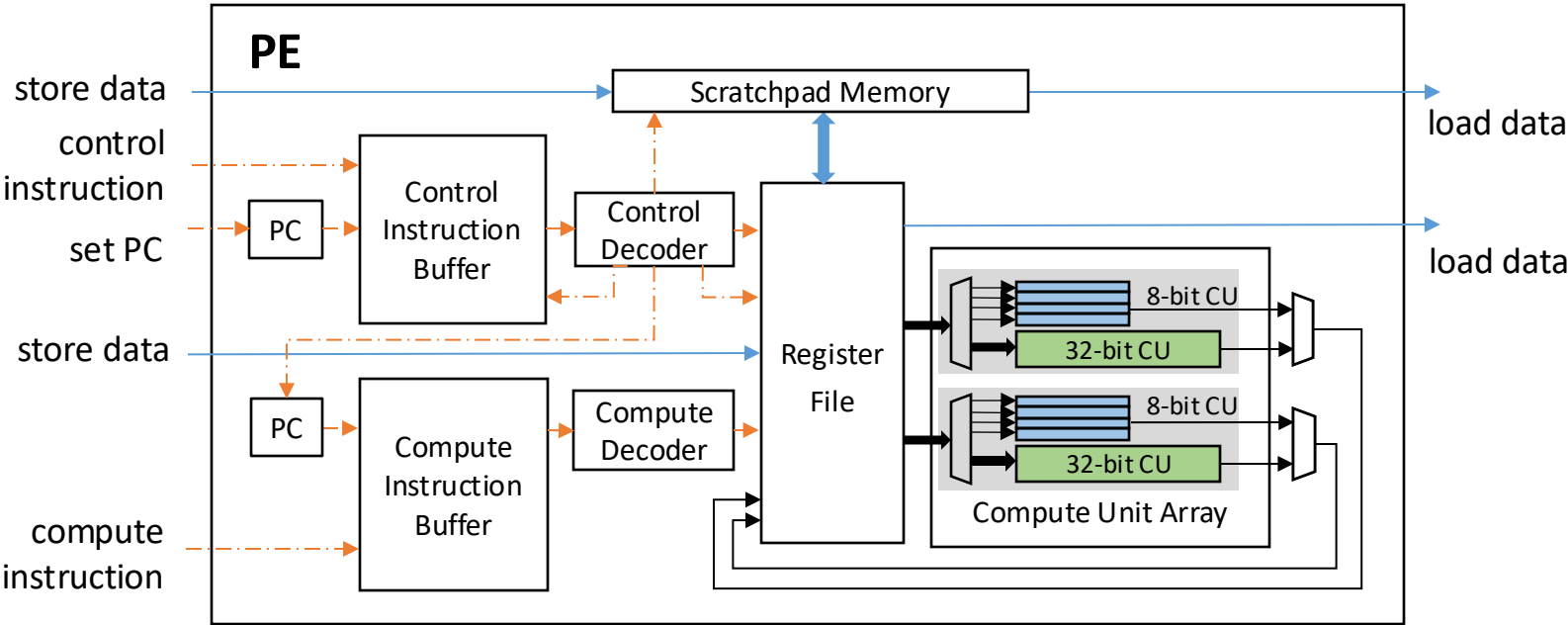
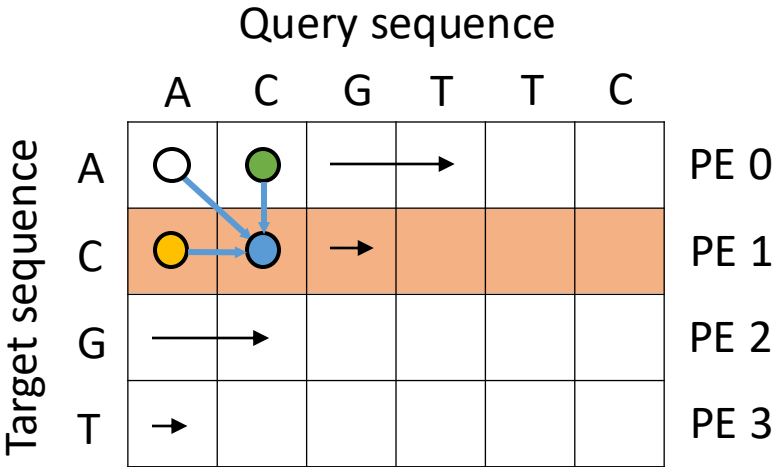


GenDP: Inter-cell dependencies and multi-precision support



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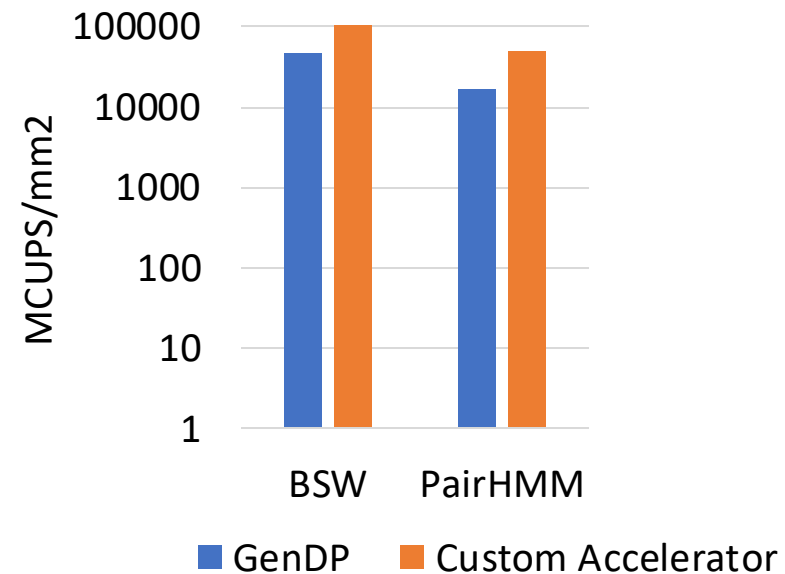
Software managed scratchpad memory for long range dependencies

GenDP performance

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- GenDP has 2.8x slowdown when compared to custom accelerators



Conclusion

- Genomic sequence data is rapidly increasing in volume and diversity
- Modern sequence analysis pipelines present unique computation challenges that multi-core CPUs and GPUs have struggled to meet
- The GenomicsBench benchmark suite is an initial attempt to capture this diversity with the goal of enabling the development of custom hardware architectures
- GenomicsBench has already helped to develop several hardware accelerators
 - Enumerated Radix Tree accelerator for memory-intensive exact match search
 - GenDP for dynamic programming acceleration
 - Pruning-based accelerator for PairHMM

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David Patterson

Intel Labs

Sanchit Misra

Md. Vasimuddin

Somnath Paul

illumina[®]

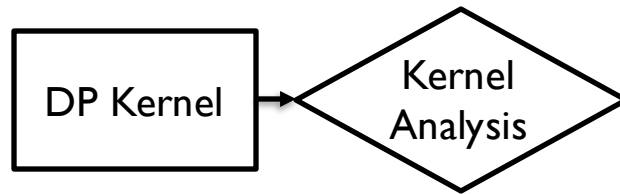


GenDP: A General Dynamic Programming Accelerator for Genomics

- **DPAx**: programmable dynamic programming (DP) accelerator.
- **DPMaP**: map the objective function of DP algorithm to DPAx accelerator.

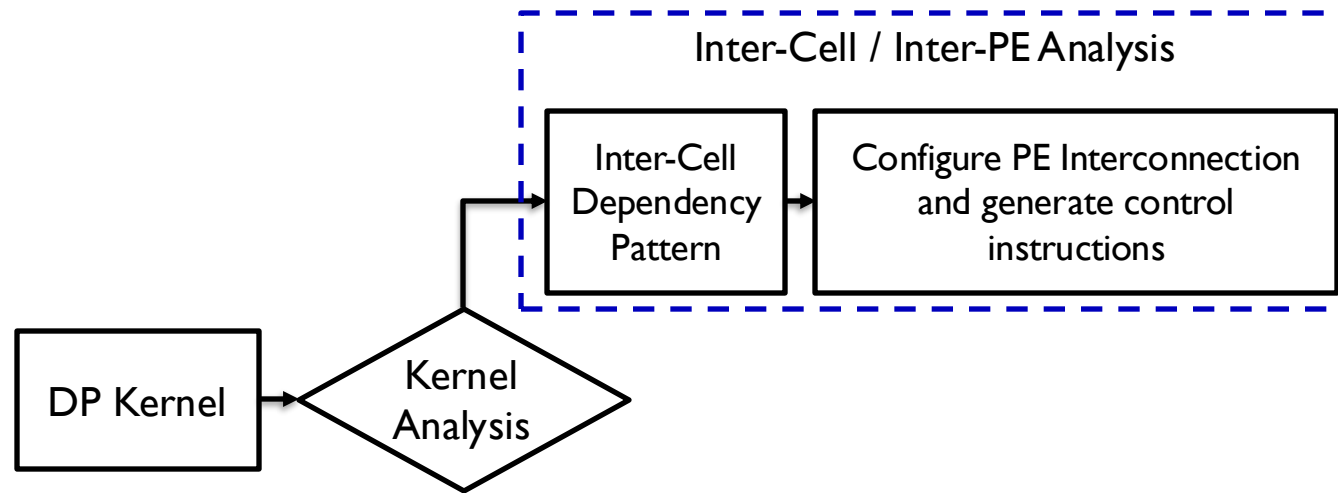
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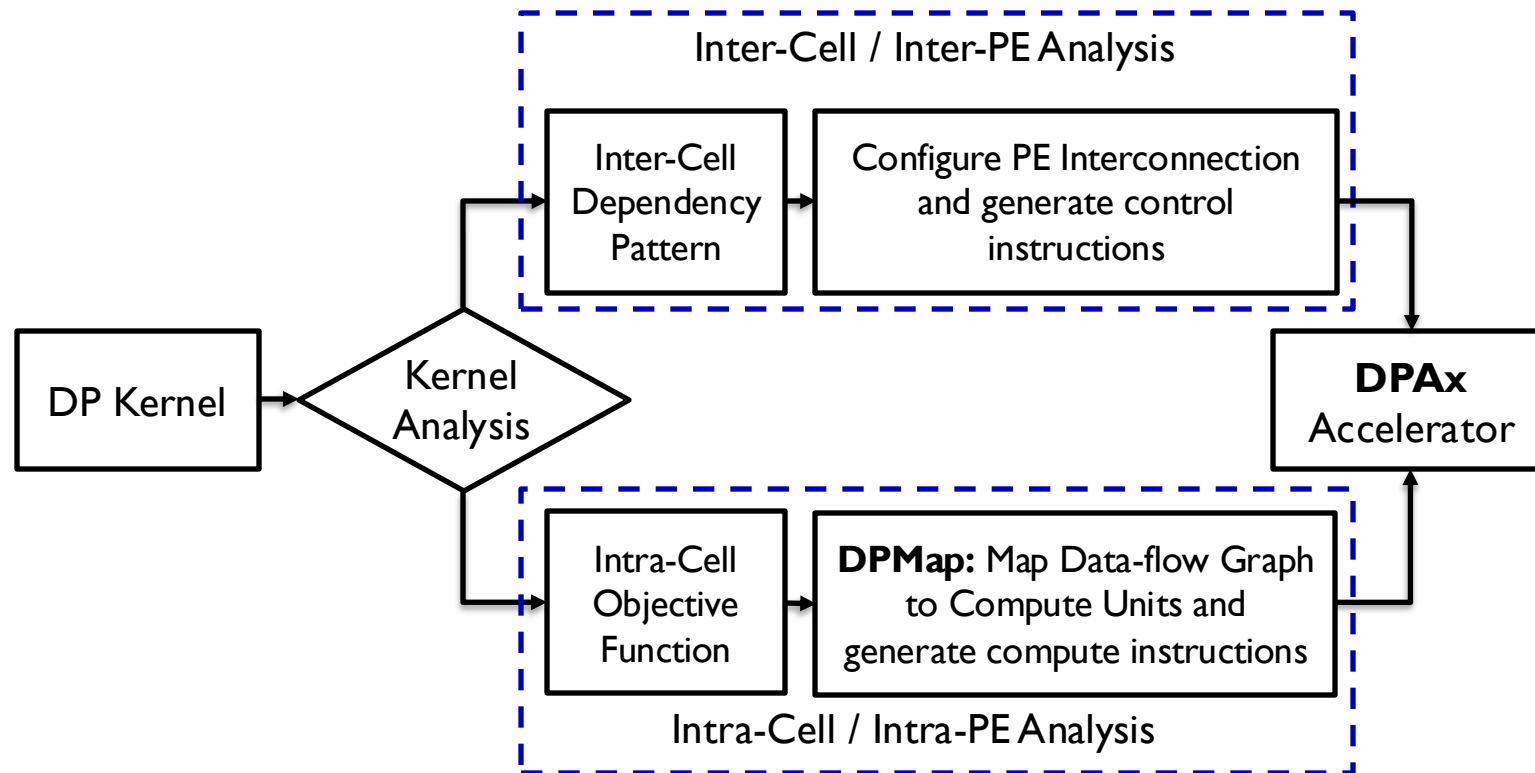
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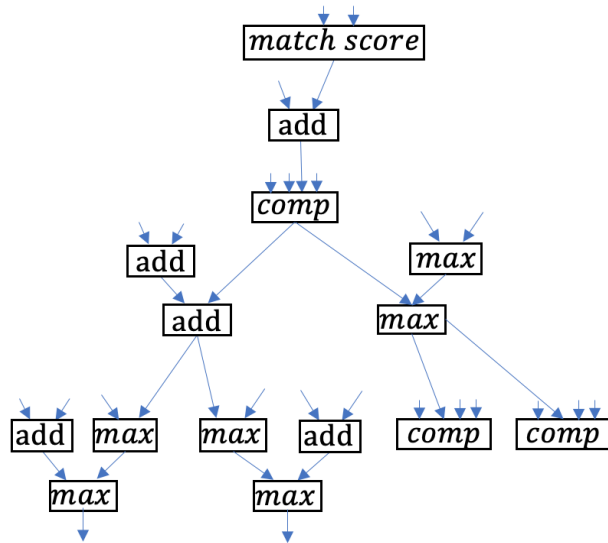


DPMap Algorithm

- **DPMap**: map the objective function of DP algorithm to programmable compute units in DPAx.

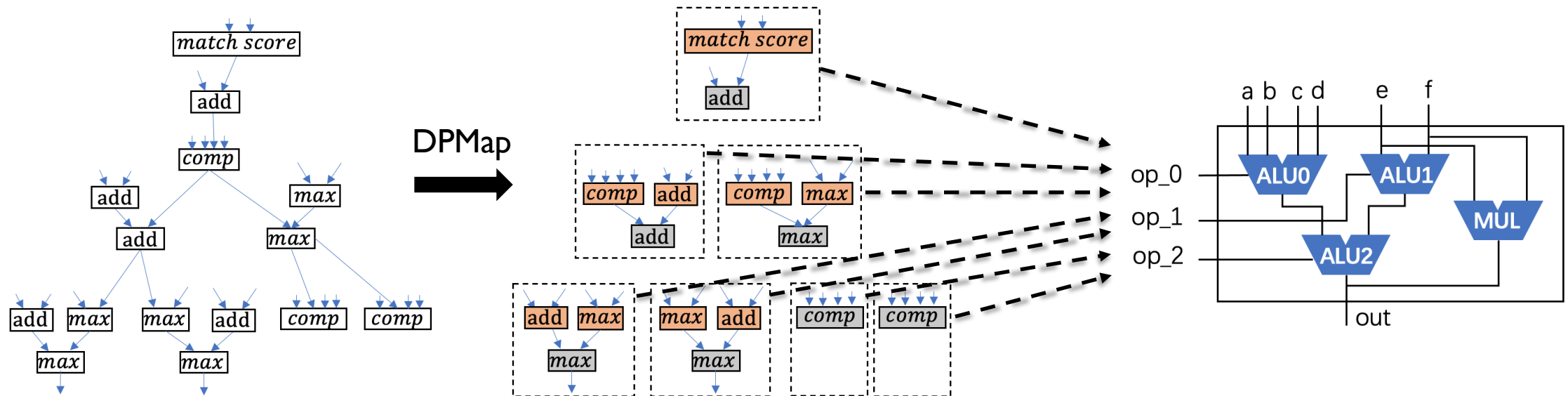
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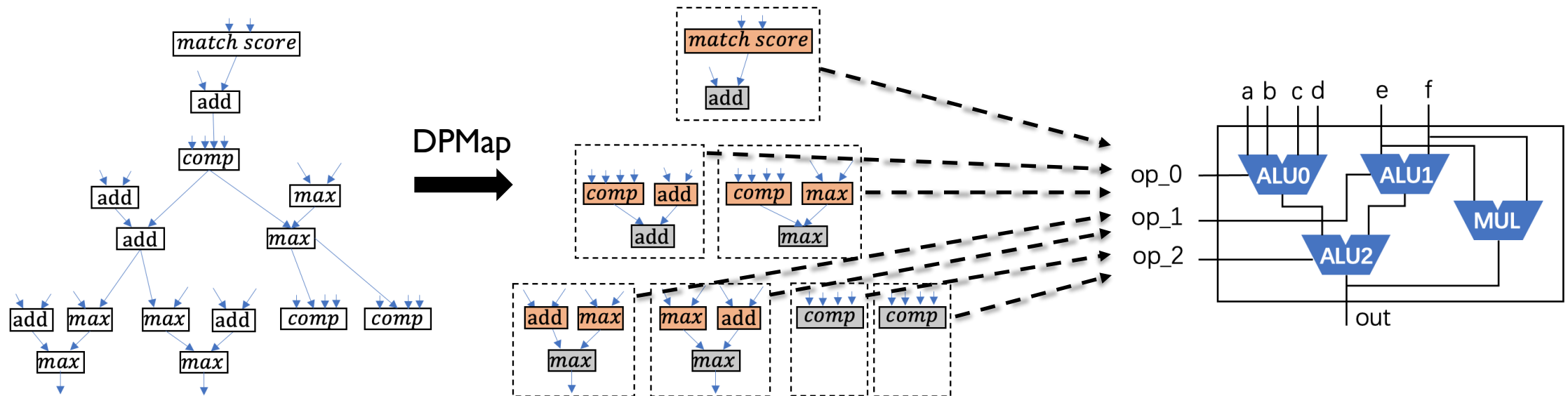
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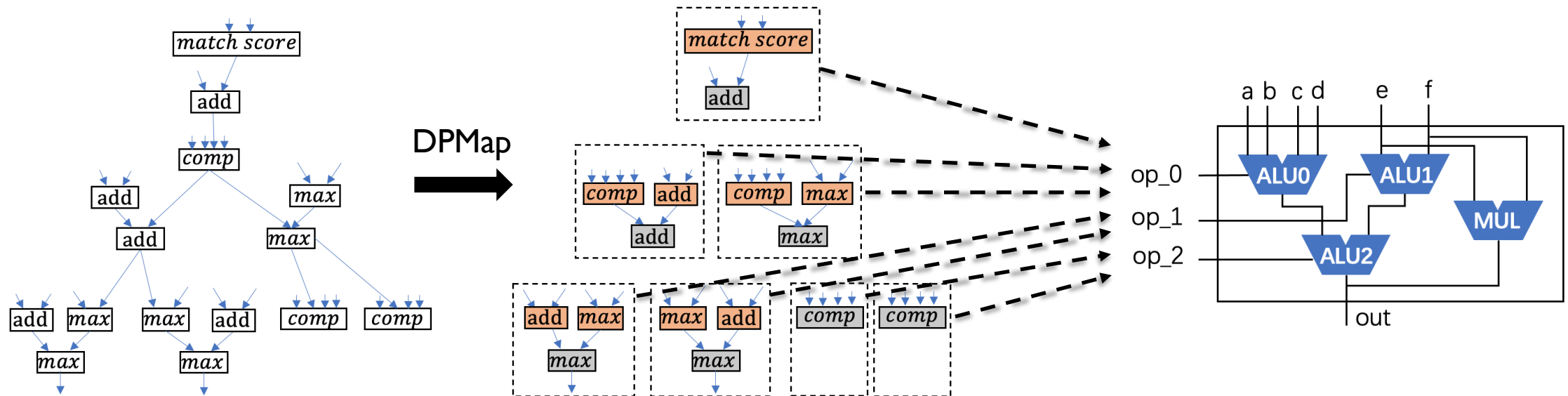
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 - **Refinement**: Break the single-strand structure



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More details
in the paper

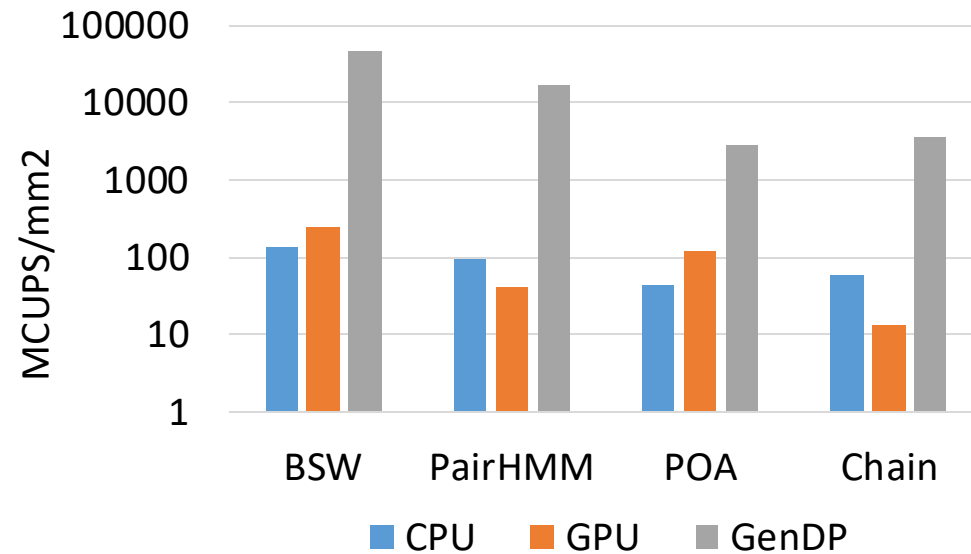


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